ค่าถำคับชั้นน้อยที่สุดของกราฟ

นายศราวุฒิ รัตนประยูร

# ศูนย์วิทยทรัพยากร จุฬาลงกรณ์มหาวิทยาลัย

วิทยานิพนธ์นี้เป็นส่วนหนึ่งของการศึกษาตามหลักสูตรปริญญาวิทยาศาสตรมหาบัณฑิต สาขาวิชาคณิตศาสตร์ ภาควิชาคณิตศาสตร์ คณะวิทยาศาสตร์ จุฬาลงกรณ์มหาวิทยาลัย ปีการศึกษา 2552 ลิขสิทธิ์ของจุฬาลงกรณ์มหาวิทยาลัย

#### MINIMUM RANK OF GRAPHS

Mr. Sarawut Rattanaprayoon

## ู่เนย่วิทยทรัพยากร

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Thesis Title	MINIMUM RANK OF GRAPHS
Ву	Mr. Sarawut Rattanaprayoon
Field of Study	Mathematics
Thesis Advisor	Associate Professor Wanida Hemakul, Ph.D.
Thesis Co-Advisor	Thiradet Jiarasuksakun, Ph.D.

Accepted by the Faculty of Science, Chulalongkorn University in

Partial Fulfillment of the Requirements for the Master's Degree

S. Harmanghera Dean of the Faculty of Science

(Professor Supot Hannongbua, Dr.rer.nat.)

THESIS COMMITTEE

Inchit Termutten Chairman

(Associate Professor Imchit Termwuttipong, Ph.D.)

W. Hemabul ...... Thesis Advisor

(Associate Professor Wanida Hemakul, Ph.D.)

Flint filh ..... Thesis Co-Advisor

(Thiradet Jiarasuksakun, Ph.D.)

When heal . Examiner

(Associate Professor Wicharn Lewkeeratiyutkul, Ph.D.)

J. Suliyou External Examiner

(Associate Professor Somporn Sutinuntopas, Ph.D.)

ศราวุฒิ รัตนประยูร : ค่าลำคับชั้นน้อยที่สุดของกราฟ. (MINIMUM RANK OF GRAPHS) อ. ที่ปรึกษาวิทยานิพนธ์หลัก : รศ. คร. วนิดา เหมะกุล, อ. ที่ปรึกษาวิทยานิพนธ์ ร่วม : อ. คร. ธีระเคช เจียรสุขสกุล, 56 หน้า.

ก่าถำคับชั้นน้อยที่สุดบนฟีลด์ F ของกราฟ G กือ ก่าลำดับชั้นน้อยที่สุดที่เป็นไปได้ใน บรรดาเมทริกซ์สมมาตรบนฟีลด์ F ซึ่งสมาชิกแลวที่ i หลักที่ j (i ≠ j) ไม่เป็นศูนย์ ถ้า ij เป็นเส้น เชื่อมในกราฟ G และเป็นศูนย์ ถ้า ij ไม่เป็นเส้นเชื่อมในกราฟ G เมื่อ ศูนย์ กือ เอกลักษณ์การบวก บนฟีลด์ F เมทริกซ์เหมาะที่สุดเชิงเอกภพของกราฟ G กือ เมทริกซ์สมมาตร A ที่สมาชิกทุกตัวเป็น จำนวนเต็มแต่สมาชิกที่ไม่อยู่บนแนวทแยงมุมของเมทริกซ์ A กือ จำนวน 0, 1 หรือ -1 และสำหรับ ทุกฟีลด์ F ก่าลำดับชั้นของเมทริกซ์ A เท่ากับก่าลำดับชั้นน้อยที่สุดบนฟีลด์ F ของกราฟ G ซึ่ง สมสัณฐานกับกราฟของเมทริกซ์ A เราแนะนำกราฟพัด กราฟหนังสือ กราฟดอกบัว และกราฟ สะพานแขวน และแสดงก่าลำดับชั้นน้อยที่สุดของกราฟเหล่านี้บนทุกฟีลด์ เราใช้เมทริกซ์เหมาะ ที่สุดเชิงเอกภพเพื่อแสดงว่าก่าลำดับชั้นน้อยที่สุดของกราฟเหล่านี้ไม่ขึ้นอยู่กับฟีลด์ และให้ตัวอย่าง กราฟที่มีก่าลำดับชั้นน้อยที่สุดขึ้นอยู่กับฟีลด์

# ศูนย์วิทยทรัพยากร จุฬาลงกรณ์มหาวิทยาลัย

ภาควิชาคณิตศาสตร์	
สาขาวิชาคณิตศาสตร์	
ปีการศึกษา2552	

ลายมือชื่อนิสิต Sarawut Rat	tanaprayoon
ลายมือชื่อ อ.ที่ปรึกษาวิทยานิพนธ์หลัก	W. Hemabul
ลายบือชื่อ อ ที่ปรึกษาวิทยาบิพบธ์ร่าน J	hit fill

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SARAWUT RATTANAPRAYOON : MINIMUM RANK OF GRAPHS. THESIS ADVISOR : ASSOC. PROF. WANIDA HEMAKUL, Ph.D. THESIS CO-ADVISOR : THIRADET JIARASUKSAKUN, Ph.D., 56 pp.

The minimum rank over a field F of a graph G is the smallest possible rank among all symmetric matrices over F whose (i, j)th entry  $(i \neq j)$  is nonzero whenever ij is an edge in G and is zero otherwise, where zero is the additive identity of F. A universally optimal matrix for a graph G is an integer symmetric matrix A such that every off-diagonal entry of A is 0, 1, or -1 and for all fields F, the rank of A is the minimum rank over F of G which is isomorphic to the graph of A. The fan graph, the book graph, the lotus graph and the hanging bridge graph are introduced and the minimum rank of these graphs over any field are presented. We use universally optimal matrices for these graphs to establish field independence of minimum rank. Examples verifying lack of field independence for some graphs are provided.

## จุฬาลงกรณ์มหาวิทยาลัย

Department	:	Mathematics
Field of Study	:	Mathematics
Academic Year	:	

Student's Signature : .	Sarawut
Advisor's Signature : .	W. Hemakul
Co-Advisor's Signatur	e. Slit ful

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## CHAPTER I INTRODUCTION

The minimum rank problem is, for a given graph and a field F, to determine the smallest possible rank among symmetric matrices over F whose off-diagonal pattern of zero-nonzero entries is described by the graph. Most work on minimum rank has been on the real minimum rank problem. S. Fallat and L. Hogben [5] provided a survey of known results and discussion of the motivation for the minimum rank problem. Catalogs of minimum rank and other parameters for families of graphs [7] and small graphs [8] were developed at the American Institute of Mathematics (AIM) workshop "Spectra of families of matrices described by graphs, digraphs, and sign patterns" [2] and are available on-line; these catalogs are updated routinely. The study of minimum rank over fields other than the real numbers was initiated in [3].

The minimum rank of a graph G is *field independent* if the minimum rank of G is the same for all fields. In [4], L.M. DeAlba et al. established the field independence or dependence of minimum rank for most of the families of graphs listed in the AIM on-line minimum rank graph catalog and established the minimum rank of several additional families. For almost every graph discussed that has field independent minimum rank, they exhibited a single integer matrix that over every field has the given graph and has rank in that field equal to the minimum rank over the field (what they call a universally optimal matrix described in chapter II).

Here is the outline of this thesis.

In chapter II, we recall definitions and review results of the relevant works.

In chapter III, we introduce the fan graph, the book graph, the lotus graph, and the hanging bridge graph and establish the field independence of minimum rank for these graphs by constructing universally optimal matrices.

In chapter IV, we provide examples verifying lack of field independence of minimum rank for some graphs, such as  $P_4 \vee K_2$ ,  $C_6 \vee K_4$ , the clique path KP(5, 4), and the clique-cycle path KC(5; 4).



## CHAPTER II PRELIMINARIES

We recall definitions and review the known results that are needed in our work.

A graph G means a simple undirected graph (i.e., neither loops nor multiple edges allowed). Denote by V(G) and E(G) the set of vertices and edges of G, respectively. Also, |G| denotes the number of vertices in G, and xy denotes the edge in E(G) for some  $x, y \in V(G)$ .

The adjacency matrix of a graph G, denoted  $\mathcal{A}(G) = [a_{ij}]$ , is a (0, 1)-matrix such that  $a_{ij} = 1$  if and only if  $ij \in E(G)$ .

The *degree* of vertex v in a graph G, denoted d(v), is the number of vertices adjacent to v. A *leaf* is a vertex of degree 1.

The *complement* of a graph G is the graph  $\overline{G}$  such that vertex set is V(G) and for each pair  $u, v \in V(G)$ , uv is an edge of  $\overline{G}$  if and only if uv is not an edge of G.

A graph H is a subgraph of a graph G if  $V(H) \subseteq V(G)$  and  $E(H) \subseteq E(G)$ . The subgraph G[R] of G induced by  $R \subseteq V(G)$  is the subgraph with vertex set Rand edge set  $\{ij \in E(G) : i, j \in R\}$ . The subgraph induced by  $\overline{R}$  is denoted by G - R, or in the case R is a single vertex v, by G - v.

An induced subgraph H of a graph G is a *clique* if H has an edge between every pair of vertices of H. A set of subgraphs of G, each of which is a clique and such that every edge of G is contained in at least one of these cliques, is called a *clique covering* of G. Let u and v be vertices in a graph G, a u, v-path in G is a list  $u = v_0, v_1, \ldots, v_n = v$  of vertices in V(G) such that  $v_{i-1}v_i \in E(G)$  and  $v_0, v_1, \ldots$ , and  $v_n$  are all different.

A graph G is connected if it has a u, v-path in G whenever  $u, v \in V(G)$ ; otherwise, G is disconnected.

A path is a graph  $P_n$  such that  $V(P_n) = \{v_1, v_2, ..., v_n\}$  and  $E(P_n) = \{v_i v_{i+1} : i = 1, 2, ..., n-1\}$ . A cycle is a graph  $C_n$  such that  $V(C_n) = \{v_1, v_2, ..., v_n\}$  and  $E(C_n) = \{v_i v_{i+1} : i = 1, 2, ..., n-1\} \cup \{v_n v_1\}$ . A complete graph is a graph  $K_n$  such that  $V(K_n) = \{v_1, v_2, ..., v_n\}$  and  $E(K_n) = \{v_i v_j : 1 \le i < j \le n\}$ .

A vertex v of a connected graph G is a *cut-vertex* if G - v is disconnected.

A graph G is *isomorphic to* a graph H, denoted  $G \cong H$ , if there is a bijection  $f: V(G) \to V(H)$  such that  $uv \in E(G)$  if and only if  $f(u)f(v) \in E(H)$ .

The union of graphs  $G_1, G_2, \ldots$ , and  $G_n$ , denoted  $\bigcup_{i=1}^n G_i$ , is the graph with vertex set  $\bigcup_{i=1}^n V(G_i)$  and edge set  $\bigcup_{i=1}^n E(G_i)$ . When  $V(G_i) \cap V(G_j) = \emptyset$  for all  $i \neq j$ , it is called the *disjoint union* of graphs  $G_1, G_2, \ldots$ , and  $G_n$ , denoted  $G_1 + G_2 + \cdots + G_n$ . nG denotes the disjoint union of n copies of a graph G.

The complete multipartite graph, denoted  $K_{n_1,n_2,...,n_k}$ , is the complement of  $K_{n_1} + K_{n_2} + \cdots + K_{n_k}$ . When k = 2, it is called a *complete bipartite graph*. A complete bipartite graph  $K_{1,n-1}$  is called an *n*-vertex star.

The join of graphs  $G_1$  and  $G_2$  with disjoint vertex sets  $V(G_1)$  and  $V(G_2)$  and edge sets  $E(G_1)$  and  $E(G_2)$ , denoted  $G_1 \vee G_2$ , is the union of  $G_1$  and  $G_2$  together with all the edges joining  $V(G_1)$  and  $V(G_2)$ .

Next, we give the basic definitions and the association of matrices and graphs.

Let  $S_n^F$  denote the set of  $n \times n$  symmetric matrices over a field F. For  $A = [a_{ij}] \in S_n^F$ , the graph of A, denoted  $\mathcal{G}^F(A)$ , is the graph with vertex set  $\{1, 2, \ldots, n\}$ (or  $\{v_1, v_2, \ldots, v_n\}$ ) and edge set  $\{ij : a_{ij} \neq 0, 1 \leq i < j \leq n\}$ . Note that the diagonal of A is ignored in determining  $\mathcal{G}^F(A)$ . The superscript F is used because the graph of an integer matrix may vary depending on the field in which the matrix is viewed.

#### Example 2.1. Let

$$A = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}, B = \begin{bmatrix} 0 & 1 & \sqrt{2} & 0 \\ 1 & 3.1 & -1.5 & 2 \\ \sqrt{2} & -1.5 & 1 & 1 \\ 0 & 2 & 1 & 0 \end{bmatrix}, \text{ and } C = \begin{bmatrix} 0 & 1 & 3 & 0 \\ 1 & -1 & 2 & 1 \\ 3 & 2 & 2 & 1 \\ 0 & 1 & 1 & 2 \end{bmatrix}.$$

These graphs  $\mathcal{G}^{\mathbb{R}}(A)$ ,  $\mathcal{G}^{\mathbb{R}}(B)$ , and  $\mathcal{G}^{\mathbb{R}}(C)$  are the graph G and graph  $\mathcal{G}^{\mathbb{Z}_2}(C)$  is the graph H, as shown below. Note that  $\mathcal{G}^{\mathbb{R}}(C)$  is not isomorphic to  $\mathcal{G}^{\mathbb{Z}_2}(C)$ .



The minimum rank over a field F of a graph G with n vertices is

$$\operatorname{mr}^{F}(G) = \min\{\operatorname{rank}(A) : A \in S_{n}^{F}, \mathcal{G}^{F}(A) \cong G\}.$$

In case  $F = \mathbb{R}$ , the superscript  $\mathbb{R}$  may be omitted, so we write mr(G) for  $mr^{\mathbb{R}}(G)$ and  $\mathcal{G}(A)$  for  $\mathcal{G}^{\mathbb{R}}(A)$ .

The minimum rank of a graph G is *field independent* if the minimum rank of G is the same for all fields.

Recall the result from basic linear algebra.

**Proposition 2.2.** [4] Let S be a linearly dependent set of integer vectors over  $\mathbb{Q}$ . Then for every prime number p, S is linearly dependent over  $\mathbb{Z}_p$ . If A is a square integer matrix, then for every prime p, rank $\mathbb{Z}_p(A) \leq \operatorname{rank}(A)$ , and if characteristic of a field F is 0, then rank<sup>F</sup>(A) = rank(A). **Example 2.3.** Let F be any field and G be the graph as shown below.



with

$$A = \begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$

and  $\mathcal{G}^{F}(A) \cong G$ . Note that  $\operatorname{rank}(A) = 2$ . By Proposition 2.2,  $\operatorname{rank}^{F}(A) \leq 2$ . Then  $\operatorname{mr}^{F}(G) \leq \operatorname{rank}^{F}(A) \leq 2$ . Next, show that  $\operatorname{mr}^{F}(G) \geq 2$ . Let

$$B = \begin{bmatrix} d_1 & a & b & 0 \\ a & d_2 & c & d \\ b & c & d_3 & e \\ 0 & d & e & d_4 \end{bmatrix}$$

with  $\mathcal{G}^F(B) \cong G$  where  $a, b, c, d, e, d_1, d_2, d_3, d_4 \in F$  and a, b, c, d, and e are nonzero in F. Since the third and the fourth rows of B are independent, rank $(B) \ge 2$ . Then  $\operatorname{mr}^F(G) \ge 2$ . Thus  $\operatorname{mr}^F(G) = 2$  for any field F. Therefore the minimum rank of G is field independent.

In [4], L.M. DeAlba et al. defined a universally optimal matrix to establish field independence of the minimum rank as follows. Recall that when A is an integer matrix and p is prime, A can be viewed as a matrix over  $\mathbb{Z}_p$ ; the rank of A over  $\mathbb{Z}_p$  will be denoted by rank $\mathbb{Z}_p(A)$ .

A universally optimal matrix for a graph G is an integer symmetric matrix A such that every off-diagonal entry of A is 0, 1, or -1 and  $\mathcal{G}(A) \cong G$  and for all fields F, rank<sup>F</sup>(A) = mr<sup>F</sup>(G). **Example 2.4.** From example 2.3, the graph  $G \cong \mathcal{G}(A)$  where

	Г			٦
A =	0	1	1	0
	1	1	1	1
	1	1	1	1
	0	1	1	0

and rank<sup>F</sup>(A) = 2 = mr<sup>F</sup>(G) for any field F. Therefore A is a universally optimal matrix for G.

In [4], L.M. DeAlba et al. showed the results about field independence of the minimum rank for families of graphs and these graphs have universally optimal matrices which is presented in Table 2.1. Definitions of graphs in this table can be found in the Appendix.

G	$\operatorname{mr}^{F}(G)$	G	$\operatorname{mr}^F(G)$
$P_n$ (path)	n-1	$P_s \Box P_s$	$s^2 - s$
$C_n$ (cycle)	n-2	$C_s \Box C_s$	$s^2 - (s + 2 \lfloor \frac{s}{2} \rfloor)$
$K_n$ (complete graph)	1	$K_s \Box K_s$	2s - 2
$K_{p,q}$ (complete bipartite	2	claw-free block-clique	# of blocks
graph)		(i.e., line graph of tree)	
$N_s$ (necklace)	3s - 2	$K_t \circ K_s$	t+1
$P_{m,k}$ (pineapple),	3	$C_t \circ K_1, t \ge 4$	$2t - \lfloor \frac{t}{2} \rfloor$
$m \ge 3, k \ge 2$	รัญใ	หาวิทยาล	181
T (tree)		$C_t \circ K_s, s \ge 2$	2t - 2
unicyclic		$T_n$ (supertriangle)	$\frac{1}{2}n(n-1)$
polygonal path	n-2		

Table 2.1: Summary of field independence of the minimum rank over any field F for families of graphs

In [3], W. Barrett et al. showed that if

$$J = \begin{bmatrix} 1 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 1 & 1 \end{bmatrix} \text{ and } A = \begin{bmatrix} \mathbf{0}_{3\times3} & J & J \\ J & \mathbf{0}_{3\times3} & J \\ J & J & \mathbf{0}_{3\times3} \end{bmatrix}$$

where  $\mathbf{0}_{3\times 3}$  is the 3 × 3 zero matrix, then the matrix A is a universally optimal matrix for the complete multipartite graph  $K_{3,3,3}$  shown in Figure 2.1 because if characteristic of a field F is 2, rank<sup>F</sup>(A) = 2 = mr<sup>F</sup>( $K_{3,3,3}$ ); otherwise, rank<sup>F</sup>(A) = 3 = mr<sup>F</sup>( $K_{3,3,3}$ ). But  $K_{3,3,3}$  does not have field independent minimum rank.

In [4], L.M. DeAlba et al. showed that if G is the disjoint union of  $K_{3,3,3}$  and  $\overline{P_3 \cup 2K_3}$  shown in Figure 2.1, then G has field independent minimum rank but G does not have a universally optimal matrix.



Figure 2.1: The complete multipartite graph  $K_{3,3,3}$  and the graph  $\overline{P_3 \cup 2K_3}$ 

**Remark 2.5.** [4] The existence of a universally optimal matrix for the graph G implies  $\operatorname{mr}^{F}(G) \leq \operatorname{mr}(G)$  for all fields F, or equivalently, the existence of a field F such that  $\operatorname{mr}^{F}(G) > \operatorname{mr}(G)$  implies that G does not have a universally optimal matrix.

In [4], L.M. DeAlba et al. showed the results about the minimum rank of graphs are dependent of the field and these graphs does not have a universally optimal matrix which is presented in Table 2.2. Definitions of graphs in this table can be found in the Appendix.

G	$\operatorname{mr}(G)$	$\operatorname{mr}^{\mathbb{Z}_2}(G)$
$W_6$ (wheel)	3	4
$M_5$ (Möbius ladder)	6	8
$L(K_7)$	5	6
$H_3$ (half-graph)	3	4
K <sub>2,2,2,2</sub>	2	4
complement of 2-tree $H$ in Figure 2.2	4	5
complement of tree $T$ in Figure 2.3	3	4
$\overline{3K_2 \cup K_1}$	2	4
$\overline{C_6} \cong K_3 \Box K_2 \cong K_3 \Box P_2 \cong C_3 \Box P_2$	3	4
$C_5 \Box K_3 \cong C_5 \Box C_3$	9	10
$P_3 \boxtimes P_3$	4	6

Table 2.2: Summary of field dependence of the minimum rank for graphs



Figure 2.3: A tree T and its complement  $\overline{T}$ 

In chapter IV, we present some graphs which do not have a universally optimal matrix by using Remark 2.5.

We introduce the following notation about specific matrices and a vector which will be used to determine universally optimal matrices.

- 1.  $I_n$  denotes the  $n \times n$  identity matrix.
- 2.  $\mathbf{0}_{m \times n}$  denotes the  $m \times n$  zero matrix.
- 3. diag $(a_1, a_2, \ldots, a_n)$  denotes the  $n \times n$  matrix of the form

$a_1$	0	0	0	
0	$a_2$	0	0	
0	0	•••	0	Ċ
0	0	0	$a_n$	

4. diag' $(a_1, a_2, \ldots, a_{n-1})$  denotes the  $n \times n$  matrix of the form

$$\begin{bmatrix} 0 & a_1 & 0 & 0 & 0 \\ a_1 & 0 & a_2 & 0 & 0 \\ 0 & a_2 & 0 & \ddots & 0 \\ 0 & 0 & \ddots & \ddots & a_{n-1} \\ 0 & 0 & 0 & a_{n-1} & 0 \end{bmatrix}$$

5. diag<sup>"</sup> $(a_1, a_2, \ldots, a_{n-2})$  denotes the  $n \times n$  matrix of the form

$$\begin{bmatrix} 0 & 0 & a_1 & 0 & 0 & 0 \\ 0 & 0 & 0 & a_2 & 0 & 0 \\ a_1 & 0 & 0 & 0 & \ddots & 0 \\ 0 & a_2 & 0 & 0 & \ddots & a_{n-2} \\ 0 & 0 & \ddots & \ddots & \ddots & 0 \\ 0 & 0 & 0 & a_{n-2} & 0 & 0 \end{bmatrix}.$$

6. "repeat[]" means the sequence enclosed in parentheses appears as many times as needed (possibly zero times) to obtain a vector of the correct length.

For example,  $(1, 1, -1, 0, 0, -1, -1, 0, 0, -1, \dots, -1, 0, 0, -1)^T = (1, 1, \text{repeat}$  $[-1, 0, 0, -1])^T$ .

The result from the following proposition will be used to determine minimum ranks of graphs and universally optimal matrices.

#### **Proposition 2.6.** [4, 6]

1. The path  $P_n$  has a universally optimal matrix of the form  $\mathcal{A}(P_n) + D$  where

$$D = \begin{cases} \operatorname{diag}(\operatorname{repeat}[0]) & \text{if } n \text{ is odd,} \\ \\ \operatorname{diag}(\operatorname{repeat}[0], 1, 1) & \text{if } n \text{ is even.} \end{cases}$$

2. The cycle  $C_n$  has a universally optimal matrix of the form  $\mathcal{A}(C_n) + D$  where

$$D = \begin{cases} \text{diag}(\text{repeat}[0]) & \text{if } n \equiv 0 \pmod{4}, \\ \text{diag}(\underbrace{1, 1, \dots, 1}_{9}, \text{repeat}[0]) & \text{if } n \equiv 1 \pmod{4} \text{ and } n \neq 5, \\ \text{diag}(1, 1, 1, 1, 1, 1, 1, \text{repeat}[0]) & \text{if } n \equiv 2 \pmod{4}, \\ \text{diag}(1, 1, 1, 1, 1, 1, \text{repeat}[0]) & \text{if } n \equiv 3 \pmod{4}, \\ \text{diag}(0, 0, -1, -1, -1) & \text{if } n = 5. \end{cases}$$

3. The complete graph  $K_n$  has a universally optimal matrix of the form  $\mathcal{A}(K_n)$ +

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Example 2.7.

are universally optimal matrices for  $P_3$ ,  $C_4$ , and  $K_5$ , respectively.

The next results are tools to determine lower bounds or upper bounds for the minimum rank of graphs.

#### **Proposition 2.8.** [3, 5]

- 1. If H is an induced subgraph of a graph G, then  $mr^F(H) \le mr^F(G)$  for any field F.
- 2. If  $G_1, G_2, \ldots$ , and  $G_n$  are graphs and  $G = \bigcup_{i=1}^n G_i$ , then  $\operatorname{mr}(G) \leq \sum_{i=1}^n \operatorname{mr}(G_i)$ .

**Example 2.9.** We determine a lower bound and an upper bound for the minimum rank of a graph *G*.



Let F be a field. Since the path  $P_4$  is an induced subgraph of G and by Proposition 2.8 (1),  $\operatorname{mr}^F(P_4) \leq \operatorname{mr}^F(G)$ . By Table 2.1,  $\operatorname{mr}^F(P_4) = 3$ . Thus  $3 \leq \operatorname{mr}^F(G)$  for any field F. We can view that G is the union of  $K_2$  and 2 copies of  $K_3$ . By Table 2.1 and Proposition 2.8 (2),  $\operatorname{mr}(G) \leq \operatorname{mr}(K_2) + 2\operatorname{mr}(K_3) = 1 + 2 = 3$ . Thus  $\operatorname{mr}(G) \leq 3$ .

In [1], F. Barioli et al. used the idea of covering the edges of a graph with subgraphs to determine the upper bound for the minimum rank of a graph G.

An (edge) covering of a graph G is a set of subgraphs  $C = \{G_i, i = 1, 2, ..., n\}$ such that G is the union  $G = \bigcup_{i=1}^{n} G_i$ . A graph has many possible coverings, but some, such as clique coverings, are more useful than others. For a given covering C,  $c_{\mathcal{C}}(e)$  denotes the number of subgraphs that have edge e as a member. **Example 2.10.** Let G be the graph shown below.



Since  $C = \{K_5, C_4\}$  is a covering of G and  $K_5$  and  $C_4$  have only one common edge  $v_4v_5$ ,  $c_{\mathcal{C}}(v_4v_5) = 2$  and  $c_{\mathcal{C}}(e) = 1$  for every edge  $e \in E(G) \setminus \{v_4v_5\}$ .

**Proposition 2.11.** [4] Let F be a field and let G be a graph. Suppose  $C = \{G_i, i = 1, 2, ..., n\}$  is a covering of G such that for each  $G_i$  there is a universally optimal matrix of the form  $\mathcal{A}(G_i) + D_i$ , where  $D_i$  is a diagonal matrix. If char F = 0 or if char F = p and  $c_C(e) \not\equiv 0 \pmod{p}$  where p is prime and for every edge  $e \in E(G)$ , then

$$\operatorname{mr}^{F}(G) \leq \sum_{i=1}^{n} \operatorname{mr}^{F}(G_{i}).$$

**Example 2.12.** Let G be the graph shown below.



By Table 2.1,  $\operatorname{mr}^{F}(P_{3}) = 2$ ,  $\operatorname{mr}^{F}(C_{4}) = 2$  and  $\operatorname{mr}^{F}(K_{5}) = 1$  for any field F. Since  $\mathcal{C} = \{P_{3}, C_{4}, K_{5}\}$  is a covering of the graph G and  $P_{3}, C_{4}$ , and  $K_{5}$  have no common edges,  $c_{\mathcal{C}}(e) = 1$  for every edge  $e \in E(G)$ . Then  $c_{\mathcal{C}}(e) \not\equiv 0 \pmod{p}$ where p is prime. By Proposition 2.11,  $\operatorname{mr}^{F}(G) \leq \operatorname{mr}^{F}(P_{3}) + \operatorname{mr}^{F}(C_{4}) + \operatorname{mr}^{F}(K_{5}) =$ 2 + 2 + 1 = 5 for any field F. Since the path  $P_{6}$  is an induced subgraph of Gand by Table 2.1 and Proposition 2.8 (1),  $5 = \operatorname{mr}^{F}(P_{6}) \leq \operatorname{mr}^{F}(G)$  for any field F. Then  $\operatorname{mr}^{F}(G) = 5$  for any field F. By Example 2.7,  $A_{1}, A_{2}$ , and  $A_{3}$  are universally optimal matrices for  $P_3, C_4$ , and  $K_5$ , respectively. Consider

$$A = \begin{bmatrix} \mathbf{0}_{3\times3} & \mathbf{0}_{3\times3} & \mathbf{0}_{3\times4} \\ \mathbf{0}_{3\times3} & A_1 & \mathbf{0}_{3\times4} \\ \mathbf{0}_{4\times3} & \mathbf{0}_{4\times3} & \mathbf{0}_{4\times4} \end{bmatrix} + \begin{bmatrix} A_2 & \mathbf{0}_{4\times6} \\ \mathbf{0}_{6\times4} & \mathbf{0}_{6\times6} \end{bmatrix} + \begin{bmatrix} \mathbf{0}_{5\times5} & \mathbf{0}_{5\times5} \\ \mathbf{0}_{5\times5} & A_3 \end{bmatrix},$$

which is

Then rank(A) = 5 and  $\mathcal{G}(A) \cong G$ . By Proposition 2.2, rank $^{F}(A) \leq \operatorname{rank}(A) = 5$ for any field F. We have  $5 = \operatorname{mr}^{F}(G) \leq \operatorname{rank}^{F}(A) \leq 5$  for any field F. Then  $\operatorname{mr}^{F}(G) = \operatorname{rank}^{F}(A)$  for any field F. Thus A is a universally optimal matrix for G and G has field independent minimum rank.

In [1], F. Barioli et al. defined a zero forcing set as a tool to determine a lower bound for the minimum rank of a graph. First, they defined the color-change rule as follows: If G is a graph with each vertex colored either white or black, u is a black vertex of G, and exactly one neighbor v of u is white, then change the color of v to black. Given a coloring of G, the *derived coloring* is the result of applying the color-change rule until no more change are possible. A zero forcing set for a graph G is a subset Z of vertices such that if initially the vertices in Z are colored black and the remaining vertices are colored white, the derived coloring of G is all black. The zero forcing number for G, denoted Z(G), is the minimum of |Z|over all zero forcing sets  $Z \subseteq V(G)$ . The parameter Z(G) is a tool to determine a lower bound for  $\operatorname{mr}^{F}(G)$ .

The next examples show zero forcing set and zero forcing numbers for some graph.

**Example 2.13.** The graph G, as shown below, has  $\{v_3, v_4\}$  as a zero forcing set by applying the color-change rule shown in steps (a)-(d) as shown in Figure 2.4 and so  $Z(G) \leq 2$ . The derived coloring of G by the only one vertex is not all black since more than one white vertices are neighbors of a black vertex. Then any set of only one vertex of G cannot be a zero forcing set for G. Thus Z(G) = 2.



**Example 2.14.** Any set of n-2 leaves of the *n*-vertex star  $K_{1,n-1}$  is a zero forcing set for  $K_{1,n-1}$  and so  $Z(K_{1,n-1}) \leq n-2$ . The derived coloring of  $K_{1,n-1}$  by any set of n-3 vertices is not all black since there are 2 or 3 vertices left which are colored white. Then any set of n-3 vertices of  $K_{1,n-1}$  cannot be a zero forcing set for  $K_{1,n-1}$ . Thus  $Z(K_{1,n-1}) = n-2$ .

**Proposition 2.15.** [3]  $Z(P_n) = 1$ ,  $Z(C_n) = 2$  and  $Z(K_n) = n - 1$ .

**Proposition 2.16.** [1] For any graph G,  $mr^F(G) \ge |G| - Z(G)$  for any field F.

The next examples, we determine a lower bound for minimum rank over a field F of some graph G.

**Example 2.17.** Consider the graph G in Example 2.13. We have Z(G) = 2. By Proposition 2.16,  $\operatorname{mr}^{F}(G) \geq |G| - Z(G) = 5 - 2 = 3$  for any field F. Thus  $\operatorname{mr}^{F}(G) \geq 3$  for any field F.

**Example 2.18.** Consider the *n*-vertex star  $K_{1,n-1}$ . By Example 2.14,  $Z(K_{1,n-1}) = n-2$ . By Proposition 2.16,  $\operatorname{mr}^{F}(K_{1,n-1}) \geq |K_{1,n-1}| - Z(K_{1,n-1}) = n - (n-2) = 2$  for any field *F*. Thus  $\operatorname{mr}^{F}(K_{1,n-1}) \geq 2$  for any field *F*.

It is not true, if H is an induced subgraph of a graph G, then  $Z(H) \ge Z(G)$ or  $Z(H) \le Z(G)$ , as shown in the next examples.

**Example 2.19.** Consider the graph G shown below with H as an induced subgraph. We obtain  $\{v_1, v_2\}$  and  $\{v_1, v_2, v_5\}$  are zero forcing sets for G and H, respectively. Thus Z(G) = 2 < 3 = Z(H).



**Example 2.20.** Since the complete graph  $K_3$  is an induced subgraph of the complete graph  $K_5$ ,  $Z(K_3) = 2 < 4 = Z(K_5)$ .

## CHAPTER III FIELD INDEPENDENCE RESULTS

In this chapter, we introduce definitions of the book graph, the fan graph, the lotus graph, the hanging bridge graph, the path-cycle graph, and the path-clique graph and establish field independence of the minimum rank for the families of these graphs. We show that these graphs have field independent minimum rank and universally optimal matrices.

First, we present the definition of the fan graph and give results about this graph.

#### 3.1 Fan Graphs

Let n be a positive integer greater than 3. The fan graph on n vertices, denoted  $F_n$ , is the graph for which  $V(F_n) = \{v_1, v_2, \dots, v_n\}$  and  $E(F_n) = \{v_i v_n : i = 1, 2, \dots, n-1\} \cup \{v_i v_{i+1} : i = 1, 2, \dots, n-2\}.$ 

**Example 3.1.** The fan graph  $F_8$  on 8 vertices is shown below.



## **Proposition 3.2.** For $n \ge 4$ , $Z(F_n) = 2$ .

Proof. We claim that  $\{v_1, v_n\}$  is a zero forcing set for  $F_n$ , and so  $Z(F_n) \leq 2$ . Assign  $v_1$  and  $v_n$  black and the other vertices white. For all  $k, 2 \leq k \leq n-1$ , we can change the color of  $v_k$  to black since  $v_k$  is the only white vertex adjacent to  $v_{k-1}$ . Now, the derived coloring of  $F_n$  is all black. Then  $\{v_1, v_n\}$  is a zero forcing set for  $F_n$ , as desired. Thus  $Z(F_n) \leq 2$ . We see that any one vertex in  $F_n$  cannot force the remaining vertices because its degree is greater than  $1, Z(F_n) \geq 2$ . Thus  $Z(F_n) = 2$ .

Next, we will show that for any field F,  $mr^F(F_n) = n - 2$  by establishing a universally optimal matrix for  $F_n$  which yields an upper bound for  $mr^F(F_n)$ .

**Theorem 3.3.** For  $n \ge 4$ , there is a diagonal matrix D such that  $\operatorname{rank}(\mathcal{A}(F_n) + D) = n - 2$ . Moreover,  $F_n$  has field independent minimum rank, and  $\mathcal{A}(F_n) + D$  is a universally optimal matrix for  $F_n$ .

*Proof.* Let D be an  $n \times n$  diagonal matrix defined by

$$D = \begin{cases} \operatorname{diag}(0, \dots, 0, \frac{n}{2}) & \text{if } n \equiv 0 \pmod{4}, \\ \operatorname{diag}(1, 0, \dots, 0, 1, \frac{n-1}{2}) & \text{if } n \equiv 1 \pmod{4}, \\ \operatorname{diag}(1, 1, 0, \dots, 0, 1, 1, \frac{n-2}{2}) & \text{if } n \equiv 2 \pmod{4}, \\ \operatorname{diag}(1, 1, 0, \dots, 0, \frac{n-1}{2}) & \text{if } n \equiv 3 \pmod{4}. \end{cases}$$

Clearly,  $\mathcal{G}(\mathcal{A}(F_n) + D) \cong F_n$ . We exhibit two independent vectors  $\vec{z_1}$  and  $\vec{z_2}$  in the kernel of  $\mathcal{A}(F_n) + D$  to show that  $\operatorname{null}(\mathcal{A}(F_n) + D) \ge 2$ . Consider the following 4 cases:

Case  $n \equiv 0 \pmod{4}$ . Then  $\vec{z_1} = (1, 0, -1, \operatorname{repeat}[0, 1, 0, -1], 0)^T$  and  $\vec{z_2} = (0, -1, -1, \operatorname{repeat}[0, 0, -1, -1], 1)^T$ .

Case  $n \equiv 1 \pmod{4}$ . Then  $\vec{z}_1 = (-1, 1, 1, \text{repeat}[-1, -1, 1, 1], -1, 0)^T$  and  $\vec{z}_2 = (-1, 0, 0, \text{repeat}[-1, -1, 0, 0], -1, 1)^T$ . Case  $n \equiv 2 \pmod{4}$ . Then  $\vec{z}_1 = (0, -1, \text{repeat}[0, 0, -1, -1], 0, 0, -1, 1)^T$  and  $\vec{z}_2 = (-1, 1, \text{repeat}[0, -1, 0, 1], 0, -1, 1, 0)^T$ . Case  $n \equiv 3 \pmod{4}$ . Then  $\vec{z}_1 = (1, -1, \text{repeat}[0, 1, 0, -1], 0)^T$  and  $\vec{z}_2 = (-1, 0, 1)^T$ . Case  $n \equiv 3 \pmod{4}$ . Then  $\vec{z}_1 = (1, -1, \text{repeat}[0, 1, 0, -1], 0)^T$  and  $\vec{z}_2 = (-1, 0, 1)^T$ .

In any case, we obtain  $\operatorname{rank}(\mathcal{A}(F_n) + D) = n - \operatorname{null}(\mathcal{A}(F_n) + D) \leq n - 2$ . Let F be any field. By Proposition 2.16 and 3.2,  $\operatorname{mr}^F(F_n) \geq |F_n| - Z(F_n) = n - 2$ . By Remark 2.2,  $\operatorname{rank}^F(\mathcal{A}(F_n) + D) \leq \operatorname{rank}(\mathcal{A}(F_n) + D) \leq n - 2$ . We have  $n - 2 \leq \operatorname{mr}^F(F_n) \leq \operatorname{rank}^F(\mathcal{A}(F_n) + D) \leq n - 2$ . Then  $\operatorname{mr}^F(F_n) = n - 2 = \operatorname{rank}^F(\mathcal{A}(F_n) + D)$ . Thus  $\operatorname{rank}(\mathcal{A}(F_n) + D) = n - 2$ . Hence  $F_n$  has field independent minimum rank, and  $\mathcal{A}(F_n) + D$  is a universally optimal matrix for  $F_n$ .

**Example 3.4.** For the fan graph  $F_8$ ,

$$\mathcal{A}(F_8) + D = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 1 & 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 & 0 & 1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 1 & 1 & 4 \end{bmatrix}$$

is a universally optimal matrix for  $F_8$  where D = diag(0, 0, 0, 0, 0, 0, 0, 4) and  $\text{mr}^F(F_8) = 6$  for any field F.

In the next section, we present the definition of the book graph and give results about this graph.

#### **3.2** Book Graphs

Let *n* be a positive integer greater than 1. The book graph on 2(n + 1)vertices, denoted  $B_n$ , is the graph for which  $V(B_n) = \{v_1, v_2, \dots, v_{2(n+1)}\}$  and  $E(B_n) = \{v_1v_{2i+1} : i = 1, 2, \dots, n\} \cup \{v_2v_{2(i+1)} : i = 1, 2, \dots, n\} \cup \{v_iv_{i+1} : i = 1, 3, 5, \dots, 2n + 1\}.$ 

**Example 3.5.** The book graph  $B_4$  on 10 vertices is shown below.



For any field F, the next result associates a lower bound for  $mr^{F}(B_{n})$ .

**Proposition 3.6.** For  $n \ge 2$ ,  $Z(B_n) \le n$ .

Proof. We claim that  $\{v_1, v_5, v_7, v_9, \ldots, v_{2n+1}\}$  is a zero forcing set for  $B_n$  and so  $Z(B_n) \leq n$ . Assign  $v_1, v_5, v_7, v_9, \ldots$ , and  $v_{2n+1}$  black and the other vertices white. For all  $k, 3 \leq k \leq n+1$ , we can change the color of  $v_{2k}$  to black since  $v_{2k}$  is the only white vertex adjacent to  $v_{2k-1}$ . That is,  $v_6, v_8, v_{10}, v_{12}, \ldots$ , and  $v_{2n+2}$  are black vertices. Then  $v_6$  can force white vertex  $v_2$  into black, also,  $v_1$  and  $v_2$  can force white vertices  $v_3$  and  $v_4$  into black, respectively. Now, the derived coloring of  $B_n$  is all black. Thus  $\{v_1, v_5, v_7, v_9, \ldots, v_{2n+1}\}$  is a zero forcing set for  $B_n$ , as desired. Hence  $Z(B_n) \leq n$ . **Theorem 3.7.** For  $n \ge 2$ , there is a diagonal matrix D such that  $\operatorname{rank}(\mathcal{A}(B_n) + D) = n + 2$ . Moreover,  $B_n$  has field independent minimum rank, and  $\mathcal{A}(B_n) + D$  is a universally optimal matrix for  $B_n$ .

Proof. Let  $D = \operatorname{diag}(0, n-2, 1, \ldots, 1)$ . Clearly,  $\mathcal{G}(\mathcal{A}(B_n) + D) \cong B_n$ . We will exhibit n independent vectors  $\vec{z_1}, \vec{z_2}, \ldots$ , and  $\vec{z_n}$  in the kernel of  $\mathcal{A}(B_n) + D$  to show that  $\operatorname{null}(\mathcal{A}(B_n) + D) \ge n$ . Then  $\vec{z_1} = (1, 1, -1, 0, \operatorname{repeat}[0, -1])^T, \vec{z_2} = (0, 0, 1, -1, -1, 1, 0, \ldots, 0)^T, \vec{z_3} = (0, 0, 0, 0, 1, -1, -1, 1, 0, \ldots, 0)^T, \ldots, \vec{z_{n-1}} = (0, \ldots, 0, 1, -1, -1, 1, 0, 0)^T$ , and  $\vec{z_n} = (0, \ldots, 0, 1, -1, -1, 1, 0, \ldots, 0)^T, \ldots, \vec{z_{n-1}} = (0, \ldots, 0, 1, -1, -1, 1, 0, 0)^T$ , and  $\vec{z_n} = (0, \ldots, 0, 1, -1, -1, 1)^T$ . We obtain  $\operatorname{rank}(\mathcal{A}(B_n) + D) = 2n + 2 - \operatorname{null}(\mathcal{A}(B_n) + D) \le 2n + 2 - n = n + 2$ . Let F be any field. By Proposition 2.16 and 3.6,  $\operatorname{mr}^F(B_n) \ge |B_n| - Z(B_n) \ge 2n + 2 - n = n + 2$ . By Remark 2.2,  $\operatorname{rank}^F(\mathcal{A}(B_n) + D) \le \operatorname{rank}(\mathcal{A}(B_n) + D) \le n + 2$ . We have  $n + 2 \le \operatorname{mr}^F(B_n) \le$  $\operatorname{rank}^F(\mathcal{A}(B_n) + D) \le n + 2$ . Then  $\operatorname{mr}^F(B_n) = n + 2 = \operatorname{rank}^F(\mathcal{A}(B_n) + D)$ . Thus  $\operatorname{rank}(\mathcal{A}(B_n) + D) = n + 2$ . Hence  $B_n$  has field independent minimum rank, and  $\mathcal{A}(B_n) + D$  is a universally optimal matrix for  $B_n$ .

**Example 3.8.** For the book graph  $B_4$ ,

is a universally optimal matrix for  $B_4$  where D = diag(0, 2, 1, 1, 1, 1, 1, 1, 1, 1, 1) and  $\text{mr}^F(B_4) = 6$  for any field F. In the next section, we present the definition of the lotus graph and give results about this graph.

#### 3.3 Lotus Graphs

Let *n* be a positive integer greater than 2. The *lotus graph* on 2*n* vertices, denoted  $Lt_n$ , is the graph for which  $V(Lt_n) = \{v_1, v_2, ..., v_{2n}\}$  and  $E(Lt_n) = \{v_iv_{i+1} : i = 1, 2, ..., 2n-1\} \cup \{v_1v_{2n}\} \cup \{v_{2i}v_{2(i+1)} : i = 1, 2, ..., n-1\} \cup \{v_2v_{2n}\}.$ 

**Example 3.9.** The lotus graph  $Lt_5$  on 10 vertices is shown below.



For any field F, the next result associates lower bound for  $mr^F(Lt_n)$ .

**Proposition 3.10.** For  $n \ge 3$ ,  $Z(Lt_n) \le n$ .

Proof. We claim that  $\{v_1, v_2, v_4, v_6, \ldots, v_{2n-2}\}$  is a zero forcing set for  $Lt_n$  and so  $Z(Lt_n) \leq n$ . Assign  $v_1, v_2, v_4, v_6, \ldots$ , and  $v_{2n-2}$  black and the other vertices white. We can change the color of  $v_{2n}$  to black since  $v_{2n}$  is the only white vertex adjacent to  $v_1$ . For  $k = 1, 2, \ldots, n-1$ , orderly, we can change the color of  $v_{2k+1}$ to black since  $v_{2k+1}$  is the only white vertex adjacent to  $v_{2k}$ , that is  $v_3, v_5, v_7, \ldots$ , and  $v_{2n-1}$  are black vertices. Now, the derived coloring of  $Lt_n$  is all black. Then  $\{v_1, v_2, v_4, v_6, \ldots, v_{2n-2}\}$  is a zero forcing set for  $Lt_n$ , as desired. Thus  $Z(Lt_n) \leq$ n. **Theorem 3.11.** For  $n \ge 3$ , there is a matrix D such that  $\operatorname{rank}(\mathcal{A}(Lt_n) + D) = n$ . Moreover,  $Lt_n$  has field independent minimum rank, and  $\mathcal{A}(Lt_n) + D$  is a universally optimal matrix for  $Lt_n$ .

*Proof.* Let D be a  $2n \times 2n$  matrix defined by

$$D = \begin{cases} \operatorname{diag}(\operatorname{repeat}[1, 0, -1, 0], 1, 2) + \operatorname{diag}'(\operatorname{repeat}[0, 0, -2, -2], -2) \\ + \operatorname{diag}''(\operatorname{repeat}[0, 0, 0, -2], 0, 0, 0, 0) & \text{if } n \text{ is odd,} \\ \operatorname{diag}(1, 0, -1, -2, -1, \operatorname{repeat}[-2, -1, -2, -1], 0, 1, 2) \\ + \operatorname{diag}'(0, 0, -2, 0, \operatorname{repeat}[-2, 0, -2, 0], -2, -2, -2) & \text{if } n \text{ is even.} \end{cases}$$

Clearly,  $\mathcal{G}(\mathcal{A}(Lt_n) + D) \cong Lt_n$ . We exhibit n independent vectors  $\vec{z}_1, \vec{z}_2, \ldots$ , and  $\vec{z}_n$  in the kernel of  $\mathcal{A}(Lt_n) + D$  to show that  $\operatorname{null}(\mathcal{A}(Lt_n) + D) \ge n$ . Then  $\vec{z}_1 = (-1, 1, 1, 0, \ldots, 0)^T, \vec{z}_2 = (0, 0, -1, 1, 1, 0, \ldots, 0)^T, \ldots, \vec{z}_{n-1} = (0, \ldots, 0, -1, 1, 1, 0)^T$ , and  $\vec{z}_n = (-1, 0, \ldots, 0, 1, 1)^T$ . In any case, we obtain  $\operatorname{rank}(\mathcal{A}(Lt_n) + D) = 2n - \operatorname{null}(\mathcal{A}(Lt_n) + D)) \le 2n - n = n$ . Let F be any field. By Proposition 2.16 and 3.10,  $\operatorname{mr}^F(Lt_n) \ge |Lt_n| - Z(Lt_n) \ge 2n - n = n$ . By Remark 2.2,  $\operatorname{rank}^F(\mathcal{A}(Lt_n) + D) \le \operatorname{rank}(\mathcal{A}(Lt_n) + D) = n$ . We have  $n \le \operatorname{mr}^F(Lt_n) \le \operatorname{rank}^F(\mathcal{A}(Lt_n) + D) \le n$ . Then  $\operatorname{mr}^F(Lt_n) = n = \operatorname{rank}^F(\mathcal{A}(Lt_n) + D)$ . Thus  $\operatorname{rank}(\mathcal{A}(Lt_n) + D) = n$ . Hence  $Lt_n$ has field independent minimum rank, and  $\mathcal{A}(Lt_n) + D$  is a universally optimal matrix for  $Lt_n$ . **Example 3.12.** For the lotus graph  $Lt_5$ ,

$$A = \begin{vmatrix} 1 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 1 & 0 & 0 & 0 & 0 & 0 & 1 \\ 0 & 1 & -1 & -1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & -1 & 0 & -1 & -1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & -1 & 1 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & -1 & -1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 1 & -1 & 0 & -1 & 1 \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & -1 & 1 & -1 \\ 1 & 1 & 0 & 0 & 0 & 0 & 0 & 1 & -1 & 2 \end{vmatrix} = \mathcal{A}(Lt_5) + D$$

is a universally optimal matrix for  $Lt_5$  where D = diag(1, 0, -1, 0, 1, 0, -1, 0, 1, 2) + diag'(0, 0, -2, -2, 0, 0, -2, -2, -2) + diag''(0, 0, 0, -2, 0, 0, 0, 0) and  $\text{mr}^F(Lt_5) = 5$  for any field F.

In the next section, we extend the definition of the path into the hanging bridge graph and we give results about this graph.

#### 3.4 Hanging Bridge Graphs

Let *n* be a positive integer greater than 1. The hanging bridge graph on 4n vertices, denoted  $Hb_n$ , is the graph contructed from a path  $P_{3n}$  by appending *n* extra vertices, with each "extra" vertex adjacent to 3 sequential path vertices. Without loss of generality, let  $v_1, v_2, \ldots$ , and  $v_{3n}$  be the vertices on path  $P_{3n}$  such that  $v_1$  and  $v_{3n}$  have degree 2 and  $v_{3n+1}, v_{3n+2}, \ldots$ , and  $v_{4n}$  be extra vertices in  $Hb_n$ .

**Example 3.13.** The hanging bridge graph  $Hb_2$  on 8 vertices is shown below.



**Proposition 3.14.** For  $n \ge 2$ ,  $Z(Hb_n) \le n + 1$ .

Proof. We claim that n extra vertices and one vertex of degree 2 form a zero forcing set for  $Hb_n$ . Let  $V(Hb_n) = \{v_1, v_2, v_3, \ldots, v_{4n}\}$ . Assign  $v_1, v_{3n+1}, v_{3n+2}, \ldots$ , and  $v_{4n}$  black and the other vertices white. Claim that  $\{v_1, v_{3n+1}, v_{3n+2}, \ldots, v_{4n}\}$  is a zero forcing set for  $Hb_n$ . For  $k = 1, 2, \ldots, 3n$ , orderly, we can change the color of  $v_{k+1}$  to black since  $v_{k+1}$  is the only white vertex adjacent to  $v_k$ , that is  $v_2, v_3, v_4, \ldots$ , and  $v_{3n}$  are black vertices. Now, the derived coloring of  $Lt_n$  is all black. Thus  $\{v_1, v_{3n+1}, v_{3n+2}, \ldots, v_{4n}\}$  is a zero forcing set for  $Hb_n$ , as desired. Hence  $Z(Hb_n) \leq n+1$ .

Next we give result about the hanging bridge graph.

**Lemma 3.15.** For  $n \ge 2$ , there exists a diagonal matrix D such that rank $(\mathcal{A}(P_{3n}) + D) = 3n - 1$ .

*Proof.* Let D be a  $3n \times 3n$  diagonal matrix defined by

$$D = \begin{cases} \operatorname{diag}(0, \operatorname{repeat}[1], 0) & \text{if } n \text{ is odd,} \\\\ \operatorname{diag}(0, 1, 0, \operatorname{repeat}[0, 1, 1]) & \text{if } n \text{ is even.} \end{cases}$$

Clearly,  $\mathcal{G}(\mathcal{A}(P_{3n}) + D) \cong P_{3n}$ . We exhibit  $\vec{z}$  in the kernel of  $\mathcal{A}(P_{3n}) + D$  to show that  $\operatorname{null}(\mathcal{A}(P_{3n}) + D) \ge 1$ . Consider the following 2 cases: Case n is odd. Then  $\vec{z} = (\text{repeat}[1, 0, -1])^T$ .

Case *n* is even. Then  $\vec{z} = (1, 0, -1, \text{repeat}[0, 1, -1])^T$ .

In any case, we obtain  $\operatorname{rank}(\mathcal{A}(P_{3n}) + D) = 3n - \operatorname{null}(\mathcal{A}(P_{3n}) + D) \leq 3n - 1$ . By Table 2.1,  $\operatorname{mr}(P_{3n}) = 3n - 1$ . We have  $3n - 1 = \operatorname{mr}(P_{3n}) \leq \operatorname{rank}(\mathcal{A}(P_{3n}) + D) \leq 3n - 1$ . Thus  $\operatorname{rank}(\mathcal{A}(P_{3n}) + D) = 3n - 1$ .

**Theorem 3.16.** There exists a diagonal matrix  $D^*$  such that  $\operatorname{rank}(\mathcal{A}(Hb_n) + D^*) = 3n - 1$  for all  $n \geq 2$ . Moreover,  $Hb_n$  has field independent minimum rank, and  $\mathcal{A}(Hb_n) + D^*$  is a universally optimal matrix for  $Hb_n$ .

Proof. Let  $D = \operatorname{diag}(d_1, d_2, d_3, \dots, d_{3n})$  be a diagonal matrix defined in the proof of Lemma 3.15 and  $\operatorname{rank}(\mathcal{A}(P_{3n}) + D) = 3n - 1$ . Define  $D^* = \operatorname{diag}(d_1^*, d_2^*, \dots, d_{4n}^*)$ where  $d_i^* = d_i$  for all  $i = 1, 2, 3, \dots, 3n$  and  $d_j^* = 1$  for all  $j = 3n + 1, 3n + 2, 3n + 3, \dots, 4n$ . Clearly,  $\mathcal{G}(\mathcal{A}(Hb_n) + D^*) \cong Hb_n$ . The matrix  $\mathcal{A}(Hb_n) + D^*$  has n duplicate rows and columns that can be deleted to leave  $\mathcal{A}(P_{3n}) + D$  without changing the rank, that is  $\operatorname{rank}(\mathcal{A}(Hb_n) + D^*) = \operatorname{rank}(\mathcal{A}(P_{3n}) + D) = 3n - 1$ . Let F be any field. By Table 2.1,  $\operatorname{mr}^F(P_{3n}) = 3n - 1$ . Since  $P_{3n}$  is an induced subgraph of  $Hb_n$  and by Proposition 2.8 (1),  $\operatorname{mr}^F(Hb_n) \ge \operatorname{mr}^F(P_{3n}) = 3n - 1$ . By Remark 2.2,  $\operatorname{rank}^F(\mathcal{A}(Hb_n) + D^*) \le \operatorname{rank}(\mathcal{A}(Hb_n) + D^*) = 3n - 1$ . We have  $3n - 1 \le \operatorname{mr}^F(Hb_n) \le \operatorname{rank}^F(\mathcal{A}(Hb_n) + D^*) \le 3n - 1$ . Then  $\operatorname{mr}^F(Hb_n) = 3n - 1 = \operatorname{rank}^F(\mathcal{A}(Hb_n) + D^*)$ . Hence  $Hb_n$  has field independent minimum rank, and  $\mathcal{A}(Hb_n) + D^*$  is a universally optimal matrix for  $Hb_n$ . □ **Example 3.17.** For the hanging bridge graph  $Hb_2$ ,

$$A = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 0 & 1 & 0 \\ 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 0 & 1 \end{bmatrix} = \mathcal{A}(Hb_n) + D$$

is a universally optimal matrix for  $Hb_2$  where D = diag(0, 1, 0, 0, 1, 1, 1, 1) and  $\text{mr}^F(Hb_2) = 5$  for any field F.

In the next example, we will construct the graph G by adding some "extra" vertex and appropriate edges to a hanging bridge graph  $Hb_n$  such that  $Hb_n$ is an induced subgraph of G and the minimum rank over a field F of G is equal to the minimum rank over F of  $Hb_n$ .

Example 3.18. Let

															1
	1	0	1	0	0	0	0	1	1	1	1	0	0	0	
	-	1	1	1	0	0	0	1	1	1	1	0	0	0	
	010	0	1	0	1	0	0	1	1	1	1	0	0	0	
	191	0	0	1	0	1	0	0	0	0	0	1	1	1	
		0	0	0	1	1	1	0	0	0	0	1	1	1	
	21	0	0	0	0	1	1	0	0	0	0	1	1	1	21
	A =	1	1	1	0	0	0	1	1	1	1	0	0	0	~
		1	1	1	0	0	0	1	1	1	1	0	0	0	
		1	1	1	0	0	0	1	1	1	1	0	0	0	
		1	1	1	0	0	0	1	1	1	1	0	0	0	
		0	0	0	1	1	1	0	0	0	0	1	1	1	
		0	0	0	1	1	1	0	0	0	0	1	1	1	
		0	0	0	1	1	1	0	0	0	0	1	1	1	
														-	

with  $\operatorname{rank}(A) = 5$  and G be the graph obtained from  $Hb_2$  by adding 5 extra

vertices and appropriate edges as shown below. We will show that  $mr^F(G) = mr^F(Hb_2)$  for any field F.



Let F be any field. By Theorem 3.16,  $\operatorname{mr}^F(Hb_2) = 5$ . Since  $Hb_2$  is an induced subgraph of G and by Proposition 2.8 (1),  $\operatorname{mr}^F(Hb_2) \leq \operatorname{mr}^F(G)$ . By Remark 2.2,  $\operatorname{rank}^F(A) \leq \operatorname{rank}(A) = 5$ . We have  $5 \leq \operatorname{mr}^F(G) \leq \operatorname{rank}^F(A) \leq 5$ . Then  $\operatorname{mr}^F(G)$  $= 5 = \operatorname{rank}^F(A)$ . Clearly,  $\mathcal{G}(A) \cong G$ . Thus A is a universally optimal matrix for G. Hence G has a universally optimal matrix, field independent minimum rank, and  $\operatorname{mr}^F(G) = \operatorname{mr}^F(Hb_2)$  for any field F.

In [4], L.M. DeAlba et al. showed that a necklace with s diamonds  $N_s$  has a universally optimal matrix  $\mathcal{A}(N_s) + I_{4s}$ , has field independent minimum rank, and  $\mathrm{mr}^F(N_s) = 3s - 2$  for any field F.

In the next example, we will construct the graph G by adding the "extra" vertex and appropriate edges to a necklace with s diamonds  $N_s$  such that  $N_s$  is an induced subgraph of G and the minimum rank over a field F of G is equal to the minimum rank over F of  $N_s$ . Example 3.19.



with rank(A) = 7 and G be the graph obtained from  $N_3$  by adding 5 extra vertices and appropriate edges as shown above. We will show that  $\operatorname{mr}^F(G) = \operatorname{mr}^F(N_3)$ for any field F. Let F be any field. We know that  $\operatorname{mr}^F(N_3) = 7$ . Since  $N_3$  is an induced subgraph of G and by Proposition 2.8 (1),  $\operatorname{mr}^F(N_3) \leq \operatorname{mr}^F(G)$ . By Remark 2.2,  $\operatorname{rank}^F(A) \leq \operatorname{rank}(A) = 7$ . We have  $7 \leq \operatorname{mr}^F(G) \leq \operatorname{rank}^F(A) \leq 7$ . Then  $\operatorname{mr}^F(G) = 7 = \operatorname{rank}^F(A)$ . Clearly,  $\mathcal{G}(A) \cong G$ . Thus A is a universally

Let

optimal matrix for G. Hence G has a universally optimal matrix, field independent minimum rank, and  $\operatorname{mr}^{F}(G) = \operatorname{mr}^{F}(N_{3})$  for any field F.

In next section, we give the definition of the path-cycle graph and show that this graph has field independent minimum rank directly. Then we determine a universally optimal matrix for this graph as well.

#### 3.5 Path-cycle Graphs

Let k be a positive integer. A path-cycle graph, denoted  $PC(m_1, m_2, \ldots, m_k;$   $n_1, n_2, \ldots, n_{k-1})$ , is obtained from paths  $P_{m_1}, P_{m_2}, \ldots$ , and  $P_{m_k}$  and cycles  $C_{n_1}, C_{n_2},$   $\ldots$ , and  $C_{n_{k-1}}$  constructed so that for  $i = 2, 3, \ldots, k$  and  $j < i, V(P_{m_{i-1}}) \cap V(C_{n_{i-1}})$ and  $V(P_{m_i}) \cap V(C_{n_{i-1}})$  have exactly one vertex and  $V(P_{m_j}) \cap V(P_{m_i}), V(C_{n_{j-1}}) \cap V(C_{n_{j-1}}) \cap V(C_{n_{i-1}})$ ,  $V(C_{n_{j-1}}) \cap V(P_{m_i})$ , and  $V(P_{m_{j-1}}) \cap V(C_{n_{i-1}})$  have no vertices. Clearly,  $|PC(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})| = \sum_{i=1}^k m_i + \sum_{i=1}^{k-1} n_i - 2(k-1).$ 

**Example 3.20.** The path-cycle graph PC(1, 3, 4; 5, 4) is shown below.

**Proposition 3.21.** For  $k \ge 1$ ,  $Z(PC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1})) \le k$ .

*Proof.* Let  $v_2, v_3, \ldots$ , and  $v_k$  be any vertex of degree 2 in  $C_{n_1}, C_{n_2}, \ldots$ , and  $C_{n_{k-1}}$ , respectively such that each is adjacent to the common vertex of  $P_{m_i}$  and  $C_{n_j}$ . If  $m_1 = 1$ , then let  $v_1$  be the common vertex of  $P_{m_1}$  and  $C_{n_1}$ ; otherwise, let  $v_1$ be the end vertex of  $P_{m_1}$  but not the common vertex of  $P_{m_1}$  and  $C_{n_1}$ . Then  $\{v_1, v_2, \ldots, v_k\}$  is a zero forcing set of  $PC(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})$  because there is only one white vertex adjacent to a black vertex so the derived coloring of  $PC(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})$  is all black. Thus  $Z(PC(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})) \leq k$ .

**Proposition 3.22.** For  $k \ge 1$ ,  $mr^F(PC(m_1, m_2, ..., m_k; n_1, n_2, ..., n_{k-1})) = \sum_{i=1}^{k} m_i + \sum_{i=1}^{k-1} n_i - 2(k-1) - k$  for any field F. Thus  $PC(m_1, m_2, ..., m_k; n_1, n_2, ..., n_{k-1})$  has field independent minimum rank.

*Proof.* Let F be any field. By Proposition 2.16 and 3.21,  $\sum_{i=1}^{k} m_i + \sum_{i=1}^{k-1} n_i - 2(k-1) - 2(k-1)$  $k \leq |PC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1})| - Z(PC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_k; n_1, n_1, n_2, \dots, n_k; n_1, n_1, \dots, n_k; n$  $(n_{k-1})) \leq \operatorname{mr}^{F}(PC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1}))).$  Let  $S = \{P_{m_1}, P_{m_2}, \dots, n_{k-1}\}$  $n_2, \ldots, n_{k-1}$ ). By Proposition 2.6, for any  $i, j, 1 \leq i \leq k$  and  $1 \leq j \leq k-1$ ,  $P_{m_i}$  and  $C_{n_j}$  have universally optimal matrices of the from  $\mathcal{A}(P_{m_i}) + D_i$  and  $\mathcal{A}(C_{n_j}) + D_j^*$ , respectively where  $D_i$  and  $D_j^*$  are diagonal matrices. Let e be any edge in  $PC(m_1, m_2, ..., m_k; n_1, n_2, ..., n_{k-1})$ . Since for i = 2, 3, ..., k and j < i,  $E(P_{m_j}) \cap E(P_{m_i}), E(C_{n_{j-1}}) \cap E(C_{n_{i-1}}), E(P_{m_i}) \cap E(C_{n_j}) \text{ and } E(P_{m_1}) \cap E(C_{n_i}) \text{ have}$ no edges,  $c_S(e) = 1$ . By Proposition 2.11,  $mr^F(PC(m_1, m_2, ..., m_k; n_1, n_2, ..., m_k; n_1, n_2, ..., m_k; n_1, n_2, ..., n_k; n_1, n_1, n_2, ..., n_k; n_1, n_1, n_2, ..., n_k; n_1, n_1, n_2, ..., n_k; n_1, n_1, n_1, n_2, ..., n_k; n_1, n_1, n_2, ..., n_k; n$  $(n_{k-1}) \leq \sum_{i=1}^{n} \operatorname{mr}^{F}(P_{m_{i}}) + \sum_{i=1}^{n-1} \operatorname{mr}^{F}(C_{n_{i}}).$  By Table 2.1, for any  $i, j, 1 \leq i \leq k$ and  $1 \leq j^{i=1} \leq k-1$ ,  $\operatorname{mr}^{F(P_{m_i})} = m_i - 1$  and  $\operatorname{mr}^{F}(C_{n_j}) = n_j - 2$ . We have  $\operatorname{mr}^{F}(PC(m_{1}, m_{2}, \dots, m_{k}; n_{1}, n_{2}, \dots, n_{k-1})) \leq \sum_{i=1}^{k} \operatorname{mr}^{F}(P_{m_{i}}) + \sum_{i=1}^{k-1} \operatorname{mr}^{F}(C_{n_{i}}) =$  $(m_1 - 1) + (m_2 - 1) + \dots + (m_k - 1) + (n_1 - 2) + (n_2 - 2) + \dots + (n_{k-1} - 2) =$  $m_1 + m_2 + \dots + m_k - k + n_1 + n_2 + \dots + n_{k-1} - 2(k-1) = \sum_{i=1}^n m_i + \sum_{i=1}^n n_i - 2(k-1) - k.$ Thus  $\operatorname{mr}^{F}(PC(m_{1}, m_{2}, \dots, m_{k}; n_{1}, n_{2}, \dots, n_{k-1})) = \sum_{i=1}^{k} m_{i} + \sum_{i=1}^{k-1} n_{i} - 2(k-1) - k$ for any field F.  We also establish a universally optimal matrix for  $PC(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})$ .

**Proposition 3.23.** For  $k \ge 1$ ,  $PC(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})$  has a universally optimal matrix.

Proof. Let  $A_1, A'_1, A_2, A'_2, \ldots, A_{k-1}, A'_{k-1}$ , and  $A_k$  be universally optimal matrices for  $P_{m_1}, C_{n_1}, P_{m_2}, C_{n_2}, \ldots, P_{m_{k-1}}, C_{n_{k-1}}$ , and  $P_{m_k}$ , respectively. Then rank $(A_i) =$  $mr(P_{m_i}) = m_i - 1$  and  $rank(A'_i) = mr(C_{n_i}) = n_i - 2$  for all i. Let  $s_j = m_1 +$  $n_1 + m_2 + n_2 + \cdots + m_{j-1} + n_{j-1} - 2(j-1) + 1$  for all  $j \in \{1, 2, \ldots, k\}$ . For  $i = 1, 2, \ldots, k$ , we construct the matrix  $B_i$  by embedding  $A_i$  into the  $s_k \times s_k$  zero matrix at the  $s_i$ th row and  $s_i$ th column as shown below



and for i = 1, 2, ..., k - 1, we define the matrix  $B'_i$  by embedding  $A'_i$  into the  $s_k \times s_k$  zero matrix at the  $(s_i + m_i - 1)$ th row and  $(s_i + m_i - 1)$ th column as shown below



We see that  $\operatorname{rank}(A_i) = \operatorname{rank}(B_i)$  and  $\operatorname{rank}(A'_i) = \operatorname{rank}(B'_i)$  for all *i*. Let  $A = \sum_{i=1}^{k-1} (B_i + B'_i) + B_k$ . Clearly,  $\mathcal{G}(A) \cong PC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1})$ . We obtain  $\operatorname{rank}(A) \le \sum_{i=1}^k \operatorname{rank}(B_i) + \sum_{i=1}^{k-1} \operatorname{rank}(B'_i) = \sum_{i=1}^k \operatorname{rank}(A_i) + \sum_{i=1}^{k-1} \operatorname{rank}(A'_i) = \sum_{i=1}^k (m_i - 1) + \sum_{i=1}^{k-1} (n_i - 2) = \sum_{i=1}^k m_i - k + \sum_{i=1}^{k-1} n_i - 2(k - 1) = \sum_{i=1}^k m_i + \sum_{i=1}^{k-1} n_i - 2(k - 1) - k$ . Let *F* be any field. By Remark 2.2,  $\operatorname{rank}^F(A) \le \operatorname{rank}(A) \le \sum_{i=1}^k m_i + \sum_{i=1}^k m_i + \sum_{i=1}^{k-1} n_i - 2(k - 1) - k$ . By Proposition 3.22,  $\sum_{i=1}^k m_i + \sum_{i=1}^{k-1} n_i - 2(k - 1) - k = \operatorname{mr}^F(PC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1})) \le \operatorname{rank}^F(A) \le \sum_{i=1}^k m_i + \sum_{i=1}^{k-1} n_i - 2(k - 1) - k$ . Then  $\operatorname{mr}^F(PC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, m_{k-1})) = \operatorname{rank}^F(A)$ . Thus *A* is a universally optimal matrix for  $PC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, m_k; n_1, n_2, \dots, m_k; n_1, n_2, \dots, m_k)$ . **Example 3.24.** By Proposition 3.23,  $mr^F(PC(1,3,4;5,4)) = 10$  for any field F with

0	1	0	0	1	0	0	0	0	0	0	0	0
1	0	1	0	0	0	0	0	0	0	0	0	0
0	1	-1	1	0	0	0	0	0	0	0	0	0
0	0	1	-1	1	0	0	0	0	0	0	0	0
1	0	0	1	-1	1	0	0	0	0	0	0	0
0	0	0	0	1	0	1	0	0	0	0	0	0
0	0	0	0	0	1	0	1	0	1	0	0	0
0	0	0	0	0	0	1	0	1	0	0	0	0
0	0	0	0	0	0	0	1	0	1	0	0	0
0	0	0	0	0	0	1	0	1	0	1	0	0
0	0	0	0	0	0	0	0	0	1	0	1	0
0	0	0	0	0	0	0	0	0	0	1	1	1
0	0	0	0	0	0	0	0	0	0	0	1	1

is a universally optimal matrix and PC(1, 3, 4; 5, 4) has field independent minimum rank.

The definition of the path-cycle graph can be extended by replacing some cycle in path-cycle graph with a polygonal path and show that the resulting graph has a universally optimal matrix and field independent minimum rank as shown in the next example.

**Example 3.25.** The graph G as shown in Figure 3.1 consists of paths  $P_2^{(1)}$ ,  $P_3^{(2)}, P_3^{(3)}$ , and  $P_3^{(4)}$ , polygonal path  $G_1$  consisted of  $C_4^{(5)}$  and  $C_6^{(6)}$ , polygonal path  $G_2$  consisted of  $C_5^{(7)}, C_4^{(8)}$ , and  $C_6^{(9)}$ , and polygonal path  $G_3$  consisted of  $C_5^{(10)}$  and  $C_4^{(11)}$  which  $V(P_2^{(1)}) \cap V(C_4^{(5)}), V(C_6^{(6)}) \cap V(P_3^{(2)}), V(P_3^{(2)}) \cap V(C_5^{(7)}), V(C_6^{(8)}) \cap V(P_3^{(3)}), V(P_3^{(3)}) \cap V(C_4^{(6)})$ , and  $V(C_4^{(7)}) \cap V(P_3^{(4)})$  have only one vertex. We show that G has a universally optimal matrix and field independent minimum rank.



Figure 3.1: The graph G with  $mr^F(G) = 27$ 

We see that  $Z(G) \leq 4$ . Let F be any field. By Proposition 2.16,  $\operatorname{mr}^{F}(G) \geq |G| - Z(G) \geq |G| - 4 = 27$ . Let  $A_1, A_2, A_3$ , and  $A_4$  be universally optimal matrices for  $P_2, P_3, P_3$ , and  $P_3$ , respectively. In Table 2.1,  $G_1, G_2$ , and  $G_3$  have universally optimal matrices, say  $A_5, A_6$ , and  $A_7$ , respectively. For all  $i = 1, 2, 3, \ldots, 7$ , let  $B_i$ be constructed (similarly to the construction in Proposition 3.23) by embedding  $A_i$  in the appropriate place in a  $27 \times 27$  matrix with  $\operatorname{rank}(A_i) = \operatorname{rank}(B_i)$ . Let  $A = \sum_{i=1}^{7} B_i$ . Then  $\operatorname{rank}(A) \leq \sum_{i=1}^{7} \operatorname{rank}(B_i) = \sum_{i=1}^{7} \operatorname{rank}(A_i) = (2-1) + (3-1) + (3-1) + (3-1) + (3-1) + (3-2) + (11-2) + (7-2) = 27$ . We obtain that  $27 \leq \operatorname{mr}^F(G) \leq \operatorname{rank}^F(A) \leq \operatorname{rank}(A) \leq 27$ . Thus  $\operatorname{mr}^F(G) = 27 = \operatorname{rank}^F(A)$ . Clearly,  $\mathcal{G}(A) \cong G$ . Hence G has a universally optimal matrix and field independent minimum rank.

In next section, we give the definition of the path-clique graph and show that this graph has field independent minimum rank directly. Also a universally optimal matrix for this graph is determined.

#### 3.6 Path-clique Graphs

Let k be a positive integer. A path-clique graph, denoted  $PK(m_1, m_2, \ldots, m_k;$  $n_1, n_2, \ldots, n_{k-1})$ , is obtained from paths  $P_{m_1}, P_{m_2}, \ldots$ , and  $P_{m_k}$  and complete graphs  $K_{n_1}, K_{n_2}, \ldots$ , and  $K_{n_{k-1}}$  constructed so that for  $i = 2, 3, \ldots, k$  and j < i,  $V(P_{m_{i-1}}) \cap V(K_{n_{i-1}})$  and  $V(P_{m_i}) \cap V(K_{n_{i-1}})$  have exactly one vertex and  $V(P_{m_j}) \cap V(P_{m_i}), V(K_{n_{j-1}}) \cap V(K_{n_{i-1}}), V(K_{n_{j-1}}) \cap V(P_{m_i})$ , and  $V(P_{m_{j-1}}) \cap V(K_{n_{i-1}})$  have no vertices.

Clearly, 
$$|PK(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1})| = \sum_{i=1}^k m_i + \sum_{i=1}^{k-1} n_i - 2(k-1).$$

**Example 3.26.** The path-clique graph PK(1, 3, 4; 5, 4) is shown below.

**Proposition 3.27.** For  $k \ge 1$ ,  $Z(PK(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1})) \le \sum_{i=1}^{k-1} n_i$ -2k-3.

 $\begin{array}{l} Proof. \ \text{Let} \ v_1^{(j)}, v_2^{(j)}, \dots, \ \text{and} \ v_{n_j-2}^{(j)} \ \text{be any vertex of degree} \ n_j - 1 \ \text{in} \ K_{n_j} \ \text{such that} \\ \text{each is adjacent to the common vertex of} \ P_{m_j} \ \text{and} \ K_{n_j}. \ \text{If} \ m_1 = 1, \ \text{then let} \ v_0 \ \text{be the} \\ \text{common vertex of} \ P_{m_1} \ \text{and} \ K_{n_1}; \ \text{otherwise, let} \ v_0 \ \text{be the end vertex of} \ P_{m_1} \ \text{but not} \\ \text{the common vertex of} \ P_{m_1} \ \text{and} \ K_{n_1}. \ \text{Then} \ \{v_0, v_1^{(1)}, v_2^{(1)}, \dots, v_{n_{1-2}}^{(1)}, v_2^{(2)}, \dots, v_{n_{2-2}}^{(2)}, \dots, v_{n_{2-2}}^{(2)}, \dots, v_1^{(k-1)}, v_2^{(k-1)}, \dots, v_{n_{k-1}-2}^{(k-1)} \} \ \text{is a zero forcing set of} \ PK(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1}) \ \text{because there is only one white vertex adjacent to a black vertex so} \\ \text{the derived coloring of} \ PK(m_1, m_2, \dots, m_k; n_1, n_2, \dots, m_{k-1}) \ \text{is all black. Thus} \\ Z(PK(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_{k-1})) \le \sum_{i=1}^{k-1} n_i - 2k - 3. \end{array}$ 

**Proposition 3.28.** For  $k \geq 1$ ,  $mr^F(PK(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})) = \sum_{i=1}^k m_i - 1$  for any field F. Thus  $PK(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})$  has field independent minimum rank.

*Proof.* Let F be any field. By Table 2.1,  $\operatorname{mr}^{F}(P_{m_1+m_2+\cdots+m_k}) = \sum_{i=1}^{k} m_i - 1$ . Since

$$\begin{split} &P_{m_{1}+m_{2}+\dots+m_{k}} \text{ is an induced subgraph of } PK(m_{1},m_{2},\dots,m_{k};n_{1},n_{2},\dots,n_{k-1}) \\ &\text{and by Proposition 2.8 (1), } \sum_{i=1}^{k} m_{i}-1 = \operatorname{mr}^{F}(P_{m_{1}+m_{2}+\dots+m_{k}}) \leq \operatorname{mr}^{F}(PK(m_{1},m_{2},\dots,m_{k};n_{1},n_{2},\dots,n_{k-1})) \\ &\dots,m_{k};n_{1},n_{2},\dots,n_{k-1})). \text{ Let } C = \{P_{m_{1}},P_{m_{2}},\dots,P_{m_{k}},K_{n_{1}},K_{n_{2}},\dots,K_{n_{k-1}}\} \text{ and } \\ &\text{cleary } C \text{ is a covering of } PK(m_{1},m_{2},\dots,m_{k};n_{1},n_{2},\dots,n_{k-1}). \text{ By Proposition } \\ &2.6, \text{ for any } i,j,1 \leq i \leq k \text{ and } 1 \leq j \leq k-1, P_{m_{i}} \text{ and } K_{n_{j}} \text{ have universally optimal } \\ &\text{matrices of the from } \mathcal{A}(P_{m_{i}}) + D_{i} \text{ and } \mathcal{A}(K_{n_{j}}) + D_{j}^{*}, \text{ respectively where } D_{i} \text{ and } \\ &D_{j}^{*} \text{ are diagonal matrices. Let } e \text{ be any edge in } PK(m_{1},m_{2},\dots,m_{k};n_{1},n_{2},\dots,m_{k};n_{1},n_{2},\dots,n_{k-1}). \\ &E(P_{m_{i}}) \cap E(K_{n_{j}}), \text{ and } E(P_{m_{1}}) \cap E(K_{n_{i}}) \text{ have no edges, } c_{C}(e) = 1. \\ &\text{By Proposition } \\ &2.11, \operatorname{mr}^{F}(PK(m_{1},m_{2},\dots,m_{k};n_{1},n_{2},\dots,n_{k-1})) \leq \sum_{i=1}^{k} \operatorname{mr}^{F}(P_{m_{i}}) + \sum_{i=1}^{k-1} \operatorname{mr}^{F}(K_{n_{i}}) \\ &= \sum_{i=1}^{k} m_{i} - k + (k-1) = \sum_{i=1}^{k} m_{i} - 1. \\ &\text{Thus } \operatorname{mr}^{F}(PK(m_{1},m_{2},\dots,m_{k};n_{1},n_{2}$$

**Proposition 3.29.** For  $k \ge 1$ ,  $PK(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})$  has a universally optimal matrix.

Proof. Let  $A_1, A'_1, A_2, A'_2, \ldots, A_{k-1}, A'_{k-1}$ , and  $A_k$  be universally optimal matrices for  $P_{m_1}, K_{n_1}, P_{m_2}, K_{n_2}, \ldots, P_{m_{k-1}}, K_{n_{k-1}}$ , and  $P_{m_k}$ , respectively. Then rank $(A_i) =$  $\operatorname{mr}(P_{m_i}) = m_i - 1$  and  $\operatorname{rank}(A'_i) = \operatorname{mr}(K_{n_i}) = 1$  for all i. Let  $s_j = m_1 + n_1 + m_2 +$  $n_2 + \cdots + m_{j-1} + n_{j-1} - 2(j-1) + 1$  for all  $j \in \{1, 2, \ldots, k\}$ . For  $i = 1, 2, \ldots, k$ , let  $B_i$  and  $B'_i$  be constructed (similarly to the construction in Proposition 3.23) by embedding  $A_i$  into the  $s_k \times s_k$  zero matrix at the  $s_i$ th row and  $s_i$ th column with  $\operatorname{rank}(A_i) = \operatorname{rank}(B_i)$  and  $\operatorname{rank}(A'_i) = \operatorname{rank}(B'_i)$ . Again, similar argument in Proposition 3.23 is applied. We obtain  $\operatorname{mr}^F(PC(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})) = \sum_{i=1}^k m_i - 1 = \operatorname{rank}^F(A)$ . Thus A is a universally optimal matrix for  $PK(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_{k-1})$ . **Example 3.30.** By Proposition 3.29,  $mr^{F}(PK(1,3,4;5,4)) = 7$ . for any field F

with

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1	1	1	1	1	0	0	0	0	0	0	0	0
1	1	1	1	1	0	0	0	0	0	0	0	0
1	1	1	1	1	0	0	0	0	0	0	0	0
1	1	1	1	1	0	0	0	0	0	0	0	0
1	1	1	1	1	1	0	0	0	0	0	0	0
0	0	0	0	1	0	1	0	0	0	0	0	0
0	0	0	0	0	1	1	1	1	1	0	0	0
0	0	0	0	0	0	1	1	1	1	0	0	0
0	0	0	0	0	0	1	1	1	1	0	0	0
0	0	0	0	0	0	1	1	1	1	1	0	0
0	0	0	0	0	0	0	0	0	1	0	1	0
0	0	0	0	0	0	0	0	0	0	1	1	1
0	0	0	0	0	0	0	0	0	0	0	1	1
-												

is a universally optimal matrix and PK(1,3,4;5,4) has field independent minimum rank.



## CHAPTER IV FIELD DEPENDENCE RESULTS

In our work, we also present some graphs which do not have field independence of minimum rank and these graphs do not have a universally optimal matrix.

#### 4.1 The Join of Paths and Complete Graphs

Recall that for  $t \ge 3, s \ge 2, P_t \lor K_s$  is the union of graphs  $P_t$  and  $K_s$ , with disjoint vertex sets  $V(P_t)$  and  $V(K_s)$ , and all the edges joining  $V(P_t)$  and  $V(K_s)$ .

First, we compute  $mr(P_t \vee K_s)$ .

**Proposition 4.1.** For  $t \ge 3$ ,  $s \ge 2$ ,  $mr(P_t \lor K_s) = t - 1$ .

Proof. By Table 2.1,  $\operatorname{mr}(P_t) = t - 1$ . Since  $P_t$  is an induced subgraph of  $P_t \vee K_s$ and by Proposition 2.8 (1),  $\operatorname{mr}(P_t) \leq \operatorname{mr}(P_t \vee K_s)$ . We have  $t - 1 \leq \operatorname{mr}(P_t \vee K_s)$ . We will exhibit s + 1 independent vectors  $\vec{z}_1, \vec{z}_2, \ldots$ , and  $\vec{z}_{s+1}$  in the kernel of a matrix A such that  $\mathcal{G}(A) \cong P_t \vee K_s$ . Let  $V(P_t) = \{v_1, v_2, \ldots, v_t\}$  and  $V(K_s) =$  $\{v_{t+1}, v_{t+2}, \ldots, v_{t+s}\}$ . Consider the following 4 cases: Case t = 3. Let  $A = \mathcal{A}(P_3 \vee K_s) + \operatorname{diag}(0, 0, 0, \underbrace{1, \ldots, 1}_s)$ . Then  $\vec{z}_1 = (\underbrace{0, \ldots, 0}_{s+1}, 1, 1)^T$ ,  $\vec{z}_2 = (\underbrace{0, \ldots, 0}_s, 1, -1, 0)^T$ ,  $\vec{z}_3 = (\underbrace{0, \ldots, 0}_s, 1, -1, 0, 0)^T$ ,  $\ldots, \vec{z}_{s-1} = (0, 0, 0, 1, 1, 1, 0, 0)^T$ . Case  $t \equiv 0 \pmod{4}$ . Let  $A = \mathcal{A}(P_t \vee K_s) + D$  where

$$D = \operatorname{diag}(1, 0, 0, \operatorname{repeat}[0, 0, 0], 1, \underbrace{1, \dots, 1}_{s}) + \begin{bmatrix} \mathbf{0}_{t \times t} & \mathbf{0}_{t \times s} \\ \\ \mathbf{0}_{s \times t} \begin{bmatrix} \frac{t-2}{2} & \cdots & \frac{t-2}{2} \\ \vdots & \ddots & \vdots \\ \frac{t-2}{2} & \cdots & \frac{t-2}{2} \end{bmatrix}_{s \times s} \end{bmatrix}.$$

Then  $\vec{z}_1 = (\underbrace{0, \dots, 0}_{t+s-2}, 1, -1)^T, \vec{z}_2 = (\underbrace{0, \dots, 0}_{t+s-3}, 1, -1, 0)^T, \vec{z}_3 = (\underbrace{0, \dots, 0}_{t+s-4}, 1, -1, 0, 0)^T,$  $\dots, \vec{z}_{s-1} = (\underbrace{0, \dots, 0}_{t}, 1, -1, \underbrace{0, \dots, 0}_{s-2})^T, \vec{z}_s = (\operatorname{repeat}[1, -1, -1, 1], \underbrace{0, \dots, 0}_{s})^T,$  and  $\vec{z}_{s+1} = (\operatorname{repeat}[1, 0, 0, 1], -1, \underbrace{0, \dots, 0}_{s-1})^T.$ 

Case  $t \equiv 1 \pmod{4}$ . Let  $A = \mathcal{A}(P_t \lor K_s) + D$  where

$$D = \operatorname{diag}(1, 1, 0, 1, \operatorname{repeat}[0, 0, 0, 0], 1, \underbrace{1, \dots, 1}_{s}) + \begin{bmatrix} \mathbf{0}_{t \times t} & \mathbf{0}_{t \times s} \\ \\ \mathbf{0}_{s \times t} \begin{bmatrix} \frac{t-3}{2} & \cdots & \frac{t-3}{2} \\ \vdots & \ddots & \vdots \\ \frac{t-3}{2} & \cdots & \frac{t-3}{2} \end{bmatrix}_{s \times s} \end{bmatrix}.$$

Then  $\vec{z}_1 = (\underbrace{0, \dots, 0}_{t+s-2}, 1, -1)^T, \vec{z}_2 = (\underbrace{0, \dots, 0}_{t+s-3}, 1, -1, 0)^T, \vec{z}_3 = (\underbrace{0, \dots, 0}_{t+s-4}, 1, -1, 0, 0)^T,$   $\dots, \vec{z}_{s-1} = (\underbrace{0, \dots, 0}_{t}, 1, -1, \underbrace{0, \dots, 0}_{s-2})^T, \vec{z}_s = (1, -1, 0, 1, -1, \text{repeat}[-1, 1, 1, -1],$   $\underbrace{0, \dots, 0}_{s})^T, \text{ and } \vec{z}_{s+1} = (1, 0, 0, 1, 0, \text{repeat}[0, 1, 1, 0], -1, \underbrace{0, \dots, 0}_{s-1})^T.$ Case  $t \equiv 2 \pmod{4}$ . Let  $A = \mathcal{A}(P_t \lor K_s) + D$  where

$$D = \operatorname{diag}(1, \operatorname{repeat}[0, 0, 0, 0], 1, \underbrace{1, \dots, 1}_{s}) + \begin{bmatrix} \mathbf{0}_{t \times t} & \mathbf{0}_{t \times s} \\ \\ \mathbf{0}_{s \times t} \begin{bmatrix} \frac{t-2}{2} & \cdots & \frac{t-2}{2} \\ \vdots & \ddots & \vdots \\ \frac{t-2}{2} & \cdots & \frac{t-2}{2} \end{bmatrix}_{s \times s} \end{bmatrix}.$$

Then  $\vec{z}_1 = (\underbrace{0, \dots, 0}_{t+s-2}, 1, -1)^T, \vec{z}_2 = (\underbrace{0, \dots, 0}_{t+s-3}, 1, -1, 0)^T, \vec{z}_3 = (\underbrace{0, \dots, 0}_{t+s-4}, 1, -1, 0, 0)^T,$   $\dots, \vec{z}_{s-1} = (\underbrace{0, \dots, 0}_{t}, 1, -1, \underbrace{0, \dots, 0}_{s-2})^T, \vec{z}_s = (1, \text{repeat}[-1, -1, 1, 1], -1, \underbrace{0, \dots, 0}_{s})^T,$ and  $\vec{z}_{s+1} = (1, \text{repeat}[0, 0, 1, 1], 0, -1, \underbrace{0, \dots, 0}_{s-1})^T.$  Case  $t \equiv 3 \pmod{4}$  and  $t \neq 3$ . Let  $A = \mathcal{A}(P_t \vee K_s) + D$  where

$$D = \operatorname{diag}(1, 1, 0, 1, 0, 0, \operatorname{repeat}[0, 0, 0, 0], 1, \underbrace{1, \dots, 1}_{s}) + \begin{bmatrix} \mathbf{0}_{t \times t} & \mathbf{0}_{t \times s} \\ \mathbf{0}_{s \times t} \begin{bmatrix} \frac{t-3}{2} & \dots & \frac{t-3}{2} \\ \vdots & \ddots & \vdots \\ \frac{t-3}{2} & \dots & \frac{t-3}{2} \end{bmatrix}_{s \times s} \end{bmatrix}.$$

Then 
$$\vec{z}_1 = (\underbrace{0, \dots, 0}_{t+s-2}, 1, -1)^T, \vec{z}_2 = (\underbrace{0, \dots, 0}_{t+s-3}, 1, -1, 0)^T, \vec{z}_3 = (\underbrace{0, \dots, 0}_{t+s-4}, 1, -1, 0, 0)^T,$$
  
 $\dots, \vec{z}_{s-1} = (\underbrace{0, \dots, 0}_{t}, 1, -1, \underbrace{0, \dots, 0}_{s-2})^T, \vec{z}_s = (1, -1, 0, \text{repeat}[1, -1, -1, 1], \underbrace{0, \dots, 0}_{s})^T$   
and  $\vec{z}_{s+1} = (1, 0, 0, \text{repeat}[1, 0, 0, 1], -1, \underbrace{0, \dots, 0}_{s-1})^T.$   
In any case, we obtain  $s + 1 \leq \text{null}(A)$ . Then  $\text{rank}(A) = (t + s) - \text{null}(A) \leq (t + s) - (s + 1) = t - 1$ . We have  $t - 1 \leq \text{mr}(P_t \lor K_s) \leq \text{rank}(A) \leq t - 1$ . Thus  
 $\text{mr}(P_t \lor K_s) = t - 1$ .

The next example, it is shown that  $P_4 \vee K_2$  does not have field independent minimum rank.

**Example 4.2.**  $P_4 \vee K_2 \cong P_4 \vee P_2$  does not have field independent minimum rank.



Let  $A \in S_6^{\mathbb{Z}_2}$  be such that  $\mathcal{G}^{\mathbb{Z}_2}(A) \cong P_4 \vee K_2$ . We can write

$$A = \begin{bmatrix} d_1 & 1 & 1 & 1 & 1 & 1 \\ 1 & d_2 & 1 & 1 & 1 & 1 \\ 1 & 1 & d_3 & 1 & 0 & 0 \\ 1 & 1 & 1 & d_4 & 1 & 0 \\ 1 & 1 & 0 & 1 & d_5 & 1 \\ 1 & 1 & 0 & 0 & 1 & d_6 \end{bmatrix}$$

where  $d_1, d_2, \ldots, d_6 \in \mathbb{Z}_2$ . It is easily to show that vectors  $(1, 1, d_3, 1, 0, 0), (1, 1, 1, d_4, 1, 0), and <math>(1, 1, 0, 1, d_5, 1)$  are linearly independent. Then rank $(A) \ge 3$ . Suppose that rank(A) = 3. Then  $\{(1, 1, d_3, 1, 0, 0), (1, 1, 1, d_4, 1, 0), (1, 1, 0, 1, d_5, 1)\}$  is maximal independent subset of the row vector space of A. Thus  $(1, 1, 0, 0, 1, d_6) = a \cdot (1, 1, d_3, 1, 0, 0) + b \cdot (1, 1, 1, d_4, 1, 0) + c \cdot (1, 1, 0, 1, d_5, 1)$  for some  $a, b, c \in \mathbb{Z}_2$ . We obtain  $a = 1, b = 1, c = 1, d_3 = 1, d_4 = 0, d_5 = 0, \text{ and } d_6 = 1$ . Then  $(1, d_2, 1, 1, 1, 1)$  cannot be written as a linear combination of  $(1, 1, 1, 1, 0, 0), (1, 1, 1, 0, 1, 0), and (1, 1, 0, 1, 1, 1), a \text{ contradiction. Thus rank}<math>(A) \ge 4$ . Since A is arbitrary,  $\operatorname{mr}^{\mathbb{Z}_2}(P_4 \vee K_2) \ge 4$ . Let  $B \in S_6^{\mathbb{Z}_2}$  be such that

with rank(B) = 4. Clearly,  $\mathcal{G}^{\mathbb{Z}_2}(B) \cong P_4 \vee K_2$ . Then  $\operatorname{mr}^{\mathbb{Z}_2}(P_4 \vee K_2) = 4$ . By Proposition 4.1,  $\operatorname{mr}(P_t \vee K_s) = 3$ . Thus  $\operatorname{mr}(P_4 \vee K_2) = 3 < 4 = \operatorname{mr}^{\mathbb{Z}_2}(P_4 \vee K_2)$ , i.e.,  $P_4 \vee K_2$  does not have field independent minimum rank. By Remark 2.5,  $P_4 \vee K_2$  does not have a universally optimal matrix.

#### 4.2 The Join of Cycles and Complete Graphs

Recall that for  $t \ge 3, s \ge 2, C_t \lor K_s$  is the union of graphs  $C_t$  and  $K_s$ , with disjoint vertex sets  $V(C_t)$  and  $V(K_s)$ , and all the edges joining  $V(C_t)$  and  $V(K_s)$ .

First, we compute  $mr(C_t \vee K_s)$ .

**Proposition 4.3.** For  $t \ge 3, s \ge 2$ ,  $mr(C_t \lor K_s) = t - 2$ .

*Proof.* By Table 2.1,  $\operatorname{mr}(P_{t-1}) = t - 2$ . Since  $P_{t-1}$  is an induced subgraph of  $C_t \vee K_s$  and by Proposition 2.8,  $t - 2 = \operatorname{mr}(P_{t-1}) \leq \operatorname{mr}(C_t \vee K_s)$ . We will exhibit

s+2 independent vectors  $\vec{z_1}, \vec{z_2}, \ldots$ , and  $\vec{z_{s+2}}$  in the kernel of a matrix A such that  $\mathcal{G}(A) \cong C_t \lor K_s$ . Let  $V(C_t) = \{v_1, v_2, \ldots, v_t\}$  and  $V(K_s) = \{v_{t+1}, v_{t+2}, \ldots, v_{t+s}\}$ . Consider the following 4 cases:

Case t = 3. Let  $A = \mathcal{A}(C_3 \vee K_s) + I_{s+3}$ . Then  $\vec{z}_1 = (1, \underbrace{0, \dots, 0}_{s+1}, -1)^T, \vec{z}_2 = (1, \underbrace{0, \dots, 0}_{s}, -1, 0)^T, \vec{z}_3 = (1, \underbrace{0, \dots, 0}_{s-1}, -1, 0, 0)^T, \dots$ , and  $\vec{z}_{s+2} = (1, -1, \underbrace{0, \dots, 0}_{s+1})^T$ . Case t = 5. Let

	Г								7	
	0	1	0	0	1	2	2		2	
	1	-1	1	0	0	1	1	••••	1	
	0	1	-1	1	0	1	1		1	
	0	0	1	-1	1	1	1		1	
A =	1	0	0	1	0	2	2	•••	2	
	2	1	1	1	2	7	7		7	
	2	1	1	1	2	7	7		7	
	:	:	1	:	÷	:	÷	••.	÷	
	2	1	1	1	2	7	7		7	
	L								$\dashv$ (s+5)×(s+	5)

$$D = \operatorname{diag}(\underbrace{0, \dots, 0}_{t}, \underbrace{1, \dots, 1}_{s}) + \begin{bmatrix} \mathbf{0}_{t \times t} & \mathbf{0}_{t \times s} \\ \\ \mathbf{0}_{s \times t} \begin{bmatrix} \frac{t-2}{2} & \dots & \frac{t-2}{2} \\ \vdots & \ddots & \vdots \\ \frac{t-2}{2} & \dots & \frac{t-2}{2} \end{bmatrix}_{s \times s}$$

Then  $\vec{z_1} = (\operatorname{repeat}[1, 1, 0, 0], \underbrace{0, \dots, 0}_{s-1}, -1)^T, \vec{z_2} = (\operatorname{repeat}[1, 1, 0, 0], \underbrace{0, \dots, 0}_{s-2}, -1, 0)^T,$  $\vec{z_3} = (\operatorname{repeat}[1, 1, 0, 0], \underbrace{0, \dots, 0}_{s-3}, -1, 0, 0)^T, \dots, \vec{z_s} = (\operatorname{repeat}[1, 1, 0, 0], -1, \underbrace{0, \dots, 0}_{s-1})^T,$  $\vec{z_{s+1}} = (\operatorname{repeat}[1, 0, -1, 0], \underbrace{0, \dots, 0}_{s})^T, \text{ and } \vec{z_{s+2}} = (\operatorname{repeat}[0, 1, 0, -1], \underbrace{0, \dots, 0}_{s})^T.$  Case  $t \equiv 1 \pmod{4}$  and  $t \neq 5$ . Let  $A = \mathcal{A}(C_t \vee K_s) + D$  where

$$D = \text{diag}(1, 2, 1, 1, 0, \text{repeat}[0, 0, 0, 0], 1, 1, 1, 1, \underbrace{1, \dots, 1}_{s}) + \begin{bmatrix} \mathbf{0}_{t \times t} & \mathbf{0}_{t \times s} \\ \mathbf{0}_{s \times t} & \begin{bmatrix} \frac{t-5}{2} & \cdots & \frac{t-5}{2} \\ \vdots & \ddots & \vdots \\ \frac{t-5}{2} & \cdots & \frac{t-5}{2} \end{bmatrix}_{s \times s} \end{bmatrix}.$$

$$D = \operatorname{diag}(1, 1, 1, 1, 1, 1, \operatorname{repeat}[0, 0, 0, 0], \underbrace{1, \dots, 1}_{s}) + \begin{bmatrix} \mathbf{0}_{t \times t} & \mathbf{0}_{t \times s} \\ \\ \mathbf{0}_{s \times t} \begin{bmatrix} \frac{t-4}{2} & \cdots & \frac{t-4}{2} \\ \vdots & \ddots & \vdots \\ \frac{t-4}{2} & \cdots & \frac{t-4}{2} \end{bmatrix}_{s \times s} \end{bmatrix}.$$

Then  $\vec{z}_1 = (0, 1, 0, 0, 1, 0, \text{repeat}[0, 1, 1, 0], \underbrace{0, \dots, 0}_{s-1}, -1)^T, \vec{z}_2 = (0, 1, 0, 0, 1, 0, \text{repeat}[0, 1, 1, 0], \underbrace{0, \dots, 0}_{s-1}, -1, 0, 0, 1, 0, \text{repeat}[0, 1, 1, 0], \underbrace{0, \dots, 0}_{s-3}, -1, 0, 0)^T, \vec{z}_s = (0, 1, 0, 0, 1, 0, \text{repeat}[0, 1, 1, 0], -1, \underbrace{0, \dots, 0}_{s-1})^T, \vec{z}_{s+1} = (1, -1, 0, 1, -1, 0, 0, -1), \underbrace{0, \dots, 0}_{s})^T, \text{ and } \vec{z}_{s+2} = (0, 1, -1, 0, 1, -1, 1, 0, 0, -1), \underbrace{0, \dots, 0}_{s})^T.$ 

Case  $t \equiv 3 \pmod{4}$  and  $t \neq 3$ . Let  $A = \mathcal{A}(C_t \vee K_s) + D$  where

$$D = \operatorname{diag}(1, 1, 1, \operatorname{repeat}[0, 0, 0, 0], \underbrace{1, \dots, 1}_{s}) + \begin{bmatrix} \mathbf{0}_{t \times t} & \mathbf{0}_{t \times s} \\ \\ \mathbf{0}_{s \times t} & \begin{bmatrix} \frac{t-3}{2} & \cdots & \frac{t-3}{2} \\ \vdots & \ddots & \vdots \\ \frac{t-3}{2} & \cdots & \frac{t-3}{2} \end{bmatrix}_{s \times s} \end{bmatrix}.$$

Then 
$$\vec{z}_1 = (0, 1, 0, \operatorname{repeat}[0, 1, 1, 0], \underbrace{0, \dots, 0}_{s-1}, -1)^T, \vec{z}_2 = (0, 1, 0, \operatorname{repeat}[0, 1, 1, 0], \underbrace{0, \dots, 0}_{s-1}, -1, 0, 0)^T, \dots, \vec{z}_s = (0, 1, 0, \operatorname{repeat}[0, 1, 1, 0], -1, \underbrace{0, \dots, 0}_{s-1})^T, \vec{z}_{s+1} = (1, -1, 0, \operatorname{repeat}[1, 0, -1, 0], \underbrace{0, \dots, 0}_{s})^T, \text{ and}$$
  
 $\vec{z}_{s+2} = (0, 1, -1, \operatorname{repeat}[0, 1, 0, -1], \underbrace{0, \dots, 0}_{s})^T.$   
In any case, we obtain  $s + 2 \leq \operatorname{null}(A)$ . Then  $\operatorname{rank}(A) = (t + s) - \operatorname{null}(A) \leq (t + s) - (s + 2) = t - 2$ . We have  $t - 2 \leq \operatorname{mr}(C_t \lor K_s) \leq \operatorname{rank}(A) \leq t - 2$ . Thus  
 $\operatorname{mr}(C_t \lor K_s) = t - 2$ .

The next example, it is shown that  $C_6 \vee K_4$  does not have field independent minimum rank.

**Example 4.4.**  $C_6 \vee K_4$  does not have field independent minimum rank.



Let  $A \in S_{10}^{\mathbb{Z}_2}$  be such that  $\mathcal{G}^{\mathbb{Z}_2}(A) \cong C_6 \vee K_4$ . We can write

	$d_1$	1	0	0	0	1	1	1	1	1	
	1	$d_2$	1	0	0	0	1	1	1	1	
	0	1	$d_3$	1	0	0	1	1	1	1	
	0	0	1	$d_4$	1	0	1	1	1	1	
<i>A</i> =	_ 0	0	0	1	$d_5$	1	1	1	1	1	
71 —	1	0	0	0	1	$d_6$	1	1	1	1	
	1	1	1	1	1	1	$d_7$	1	1	1	
	1	1	1	1	1	1	1	$d_8$	1	1	
	1	1	1	1	1	1	1	1	$d_9$	1	
	$\lfloor 1$	1	1	1	1	1	1	1	1	$d_{10}$	

with rank(B) = 5. Clearly,  $\mathcal{G}^{\mathbb{Z}_2}(B) \cong C_6 \vee K_4$ . Then  $\operatorname{mr}^{\mathbb{Z}_2}(C_6 \vee K_4) = 5$ . By Proposition 4.3,  $\operatorname{mr}(C_6 \vee K_4) = 4$ . Thus  $\operatorname{mr}(C_6 \vee K_4) = 4 < 5 = \operatorname{mr}^{\mathbb{Z}_2}(C_6 \vee K_4)$ , i.e.,  $C_6 \vee K_4$  does not have field independent minimum rank. By Remark 2.5,  $C_6 \vee K_4$  does not have a universally optimal matrix.

#### 4.3 Clique Paths

For  $i = 1, 2, ..., k, m_i \ge 3$ . A clique path, denoted  $KP(m_1, m_2, ..., m_k)$ , is the "path" of complete graphs built from complete graphs  $K_{m_1}, K_{m_2}, ...,$  and  $K_{m_k}$ constructed so that for i = 2, 3, ..., k and j < i - 1,  $E(K_{m_{i-1}}) \cap E(K_{m_i})$  has exactly one edge and  $V(K_{m_j}) \cap V(K_{m_i})$  has no vertices.

**Remark 4.5.**  $|KP(m_1, m_2, \dots, m_k)| = \sum_{i=1}^k m_i - 2(k-1).$ 

**Proposition 4.6.**  $mr(KP(m_1, m_2, ..., m_k)) = k.$ 

Proof. Clearly,  $P_{k+1}$  is an induced subgraph of  $KP(m_1, m_2, \ldots, m_k)$ . By Table 2.1 and Proposition 2.8 (1),  $k = mr(P_{k+1}) \leq mr(KP(m_1, m_2, \ldots, m_k))$ . We can view that  $KP(m_1, m_2, \ldots, m_k)$  is the union of the complete graphs  $K_{m_1}, K_{m_2}, \ldots$ , and  $K_{m_k}$ . By Table 2.1 and Proposition 2.8 (2),  $mr(KP(m_1, m_2, \ldots, m_k)) \leq mr(K_{m_1}) + mr(K_{m_2}) + \cdots + mr(K_{m_k}) = k$ . Thus  $mr(KP(m_1, m_2, \ldots, m_k)) = k$ .  $\Box$ 

**Example 4.7.** KP(5,4) does not have field independent minimum rank.



Let  $A \in S_7^{\mathbb{Z}_2}$  be such that  $\mathcal{G}(A) \cong KP(5,4)$ . We can write

$$A = \begin{bmatrix} d_1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & d_2 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & d_3 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & d_4 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 & d_5 & 1 & 1 \\ 0 & 0 & 0 & 1 & 1 & d_6 & 1 \\ 0 & 0 & 0 & 1 & 1 & 1 & d_7 \end{bmatrix}$$

where  $d_1, d_2, \ldots, d_7 \in \mathbb{Z}_2$ . It is easily to show that vectors  $(1, 1, d_3, 1, 1, 0, 0), (1, 1, 1, d_4, 1, 1, 1)$ , and  $(0, 0, 0, 1, 1, d_6, 1)$  are linearly independent. Then rank $(A) \ge 3$ . Since A is arbitrary,  $\operatorname{mr}^{\mathbb{Z}_2}(KP(5, 4)) \ge 3$ . Let  $B \in S_7^{\mathbb{Z}_2}$  be such that

with rank(B) = 3. Clearly,  $\mathcal{G}^{\mathbb{Z}_2}(B) = KP(5,4)$ . Then  $\operatorname{mr}^{\mathbb{Z}_2}(KP(5,4)) =$ 3. By Proposition 4.6,  $\operatorname{mr}(KP(5,4)) = 2$ . Thus  $\operatorname{mr}(KP(5,4)) = 2 < 3 =$  $\operatorname{mr}^{\mathbb{Z}_2}(KP(5,4))$ , i.e., KP(5,4) does not have field independent minimum rank. By Remark 2.5, KP(5,4) does not have a universally optimal matrix.

#### 4.4 Clique-cycle Paths

For  $i = 1, 2, ..., k, m_i \ge 3$ . A clique-cycle path, denoted  $KC(m_1, m_2, ..., m_k;$  $n_1, n_2, ..., n_k)$ , is obtained from complete graphs  $K_{m_1}, K_{m_2}, ...,$  and  $K_{m_k}$  and cycles  $C_{n_1}, C_{n_2}, ...,$  and  $C_{n_k}$  constructed so that for i = 2, 3, ..., k and j < i,  $E(K_{m_1}) \cap E(C_{n_1}), E(K_{m_i}) \cap E(C_{n_i}),$  and  $E(K_{m_i}) \cap E(C_{n_{i-1}})$  have exactly one edge and  $V(K_{m_j}) \cap V(K_{m_i}), V(C_{n_j}) \cap V(C_{n_i}), V(K_{m_j}) \cap V(C_{n_i}),$  and  $V(C_{n_{j-1}}) \cap V(K_{m_i})$ have no vertices.

**Remark 4.8.** 
$$|KC(m_1, m_2, \dots, m_k; n_1, n_2, \dots, n_k)| = \sum_{i=1}^k m_i + \sum_{i=1}^k n_i - 4k + 2.$$

**Proposition 4.9.**  $mr(KC(m_1, m_2, ..., m_k; n_1, n_2, ..., n_k)) \le \sum_{i=1}^k n_i - k.$ 

*Proof.* We can view that  $KC(m_1, m_2, ..., m_k; n_1, n_2, ..., n_k)$  is the union of complete graphs  $K_{m_1}, K_{m_2}, ...,$  and  $K_{m_k}$  and cycles  $C_{n_1}, C_{n_2}, ...,$  and  $C_{n_k}$ . By Table 2.1,  $mr(K_{m_i}) = 1$  and  $mr(C_{n_i}) = n_i - 2$  for all i = 1, 2, ..., k. By Proposition 2.8 (2),  $mr(KC(m_1, m_2, ..., m_k; n_1, n_2, ..., n_k)) \leq \sum_{i=1}^k mr(K_{m_i}) + \sum_{i=1}^k mr(C_{n_i}) = k + \sum_{i=1}^k n_i - 2k = \sum_{i=1}^k n_i - k.$ 

**Example 4.10.** KC(5;4) does not have field independent minimum rank.



Let  $A \in S_7^{\mathbb{Z}_2}$  be such that  $\mathcal{G}(A) \cong KC(5; 4)$ . We can write

$$A = \begin{bmatrix} d_1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & d_2 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & d_3 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & d_4 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & d_5 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & d_6 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & d_7 \end{bmatrix}$$

where  $d_1, d_2, \ldots, d_7 \in \mathbb{Z}_2$ . It is easily to show that vectors  $(1, 1, d_3, 1, 1, 0, 0), (1, 1, 1, 1, d_5, 1, 0), (0, 0, 0, 0, 1, d_6, 1),$  and  $(0, 0, 0, 1, 0, 1, d_7)$  are linearly independent. Then rank $(A) \ge 4$ . Since A is arbitrary,  $mr^{\mathbb{Z}_2}(KC(5; 4)) \ge 4$ . Let  $B \in S_7^{\mathbb{Z}_2}$  be such that

$$B = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 \\ 0 & 0 & 0 & 1 & 0 & 1 & 0 \end{bmatrix}$$

with rank(B) = 4. Clearly,  $\mathcal{G}^{\mathbb{Z}_2}(B) = KC(5;4)$ . Then  $\operatorname{mr}^{\mathbb{Z}_2}(KC(5;4)) = 4$ . By Proposition 4.9,  $\operatorname{mr}(KC(5;4)) \leq 3$ . Clearly,  $P_4$  is an induced subgraph of KC(5;4). By Table 2.1 and Proposition 2.8,  $3 = \operatorname{mr}(P_4) \leq \operatorname{mr}(KC(5;4))$ . Then  $\operatorname{mr}(KC(5;4)) = 3$ . Thus  $\operatorname{mr}(KC(5;4)) = 3 < 4 = \operatorname{mr}^{\mathbb{Z}_2}(KC(5;4))$ , i.e., KC(5;4) does not have field independent minimum rank. By Remark 2.5, KC(5;4) does not have a universally optimal matrix.

Question. Which values of t and k that the family  $P_t \vee K_s$ ,  $C_t \vee K_s$ ,  $KP(m_1, m_2, \ldots, m_k)$ , and  $KC(m_1, m_2, \ldots, m_k; n_1, n_2, \ldots, n_k)$  have field independent minimum rank?

ศูนย์วิทยทรัพยากร จุฬาลงกรณ์มหาวิทยาลัย

#### APPENDIX

The *necklace* with s diamonds, denoted  $N_s$ , is a graph that can be constructed from a cycle  $C_{3s}$  by appending s extra vertices, with each "extra" vertex adjacent to 3 sequential cycle vertices.

The *m*, *k*-pineapple (with  $m \ge 3, k \ge 2$ ), denoted  $P_{m,k}$ , is the graph  $K_m \cup K_{1,k}$ such that a vertex in  $V(K_m) \cap V(K_{1,k})$  is the vertex of  $K_{1,k}$  of degree *k*.



A tree is a connected graph with n vertices and n-1 edges.

A *unicyclic* is a connected graph containing exactly one cycle.



A polygonal path is a "path" of cycles built from cycles  $C_{m_1}, C_{m_2}, \ldots$ , and  $C_{m_k}$ constructed so that for  $i = 2, 3, \ldots, k$  and  $j < i-1, E(C_{m_{i-1}}) \cap E(C_{m_i})$  has exactly one edge and  $E(C_{m_j}) \cap E(C_{m_i})$  has no edges.



polygonal path built from  $C_5, C_4$  and  $C_6$ 

The Cartesian product of two graphs G and H, denoted  $G \Box H$ , is the graph with vertex set  $V(G) \times V(H)$  such that (u, v) is adjacent to (u', v') if and only if (1) u = u' and  $vv' \in E(H)$ , or (2) v = v' and  $uu' \in E(G)$ .

The strong product of two graphs G and H, denoted  $G \boxtimes H$ , is the graph with vertex set  $V(G) \times V(H)$  such that (u, v) is adjacent to (u', v') if and only if (1)  $uu' \in E(G)$  and  $vv' \in E(H)$ , or (2) u = u' and  $vv' \in E(H)$ , or (3) v = v' and  $uu' \in E(G)$ .



The corona of a graph G with a graph H, denoted  $G \circ H$ , is the graph on |G||H| + |G| vertices obtained by taking one copy of G and |G| copies of H, and joining all the vertices in the *i*th copy of H to the *i*th vertex of G.

The *n*th supertriangle, denoted  $T_n$ , is a graph G with vertex set  $V(G) = \{(i, j) : i = 1, 2, ..., n \text{ and } j = 1, 2, ..., i\}$  such that (i, j) is adjacent to (i', j') if and only if (1) |i - i'| = 1 and |j - j'| = 0, or (2) |i - i'| = 0 and |j - j'| = 1, or (3) |i - i'| = 1 and |j - j'| = 1. Clearly,  $|T_n| = \frac{1}{2}n(n+1)$ .



A block of a graph G is a maximal connected subgraph of G that has no cut-vertex. A block-clique graph is a graph in which every block is a clique.

A graph is *claw-free* if it dose not contain an induced  $K_{1,3}$ .







The *n*th wheel, denoted  $W_n$ , is the graph  $K_1 \vee C_{n-1}$ .

The sth *Möbius ladder*, denoted  $M_s$ , is obtained from  $C_s \Box P_2$  by replacing one pair of parallel cycle edges with a crossed pair.







The line graph of a graph G, denoted L(G), is the graph having vertex set E(G), with two vertices in L(G) adjacent if and only if the corresponding edges share an endpoint in G. Since we require a graph to have a nonempty set of vertices, the line graph L(G) is defined only for a graph G that has at least one edge.



The sth half-graph, denoted  $H_s$ , is the graph is constructed from (disjoint) graphs  $K_s$  and  $\overline{K_s}$ , having vertices  $u_1, u_2, \ldots$ , and  $u_s$  and  $v_{s+1}, v_{s+2}, \ldots$ , and  $v_{2s}$ , respectively, by adding all edges  $u_i u_j$  such that  $i + j \leq 2s + 1$ .



A 2-tree is a graph built from  $K_3$  by adding to it one vertex at a time adjacent to exactly a pair of existing adjacent vertices.

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## VITA

Name	Mr. Sarawut Rattanaprayoon
Date of Birth	24 November 1983
Place of Birth	Surat Thanee, Thailand
Education	B.Sc. (Mathematics), (Second Class Honors),
	Prince of Songkhla University, 2006
	Grad Dip. (Teaching), Srinakharinwirot University, 2007
Scholarship	The Promotion Project for Teacher Production in Sciences
	and Mathematics
Conference	Presented a short communication on
	• "Minimum Rank Matrices of Fan Graphs and Book
	Graphs", Annual Pure and Applied Mathematics
	Conference, 25-26 May 2009, Chulalongkorn University
	• "Field independence of the minimum rank of a graph",
	The 15th Annual Meeting in Mathematics, 10-12 March
	2010, King Mongkut's University of Technology North
	Bangkok
	Attended
	• The Annual Mathematics Conference, 27-28 March 2008,
	Chulalongkorn University
	• The International Conference on Algebra and Related
	Topics, 28-30 May 2008, Chulalongkorn University